

**Research and Development for
Space Data System Standards**

**FLEXIBLE SERIALY
CONCATENATED CONVOLUTIONAL
TURBO CODES WITH NEAR-
SHANNON BOUND PERFORMANCE
FOR TELEMETRY APPLICATIONS**

EXPERIMENTAL SPECIFICATION

CCSDS 131.2-O-1

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FOREWORD

This document describes a Serially Concatenated Convolutional turbo Coding (SCCC) scheme for telemetry applications. The flexibility, performance and proper architecture of the proposed coding scheme together with a new frame structure make the scheme suitable for achieving a significantly high spectral and power efficiency while maintaining compatibility with the existing data layer protocols.

The proposed coding scheme and its associated frame structure are specifically designed to support reconfiguration of the downlink channel (variable or adaptive coding and modulation) and to provide means for reliable synchronization at the Physical layer and the Data Link layer.

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PREFACE

This document is a CCSDS Experimental Specification. Its Experimental status indicates that it is part of a research or development effort based on prospective requirements, and as such it is not considered a Standards Track document. Experimental Recommendations are intended to demonstrate technical feasibility in anticipation of a 'hard' requirement that has not yet emerged. Experimental work may be rapidly transferred onto the Standards Track should a hard requirement emerge in the future.

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1 INTRODUCTION

1.1 BACKGROUND

In recent years, there has been an ever-growing demand for high data rate in multimedia satellite communications and earth observation missions such as COSMO and METEOSAT third generation. In order to satisfy this demand, innovative modulation and coding schemes with high spectral and energy efficiency are required to match the bandwidth- and/or power-limited conditions of satellite channels.

Further improvements in the average system throughput are envisaged by using adaptive Physical layer techniques that allow for reconfiguration of the coding and modulation scheme. The change of coding and modulation scheme is envisaged based on either estimates of channel condition or a predetermined schedule (for example, as a function of variation in the elevation angle). Such improvements require more flexibility in the Physical and MAC layer of the transmitter and receiver architecture.

This document describes a new coding/modulation scheme featuring simplicity and versatility while offering performance close to theoretical bounds. The proposed scheme provides a wide range of spectral efficiency values that can be reconfigured without interrupting the data transfer over the physical channel. The implementation of the encoder can be realized with certain degrees of parallelism without any significant increase in memory or computation requirements. This feature makes the scheme attractive in particular for on-board implementation.

1.2 PURPOSE AND SCOPE

The proposed coding scheme belongs to a family of ‘turbo-like codes’ that is structured as a serial concatenation of two convolutional codes through an interleaver (reference [B1]). The main criteria used for the code design can be highlighted as follows:

- *flexibility*, i.e., use a single mother code for all modulations and code rates;
- *simplicity*, i.e., build a code as compact and simple as possible;
- *performance*, i.e., approach Shannon capacity bounds.

The proposed SCCC scheme supports flexible design of a wide range of spectral efficiency values based on a single coding structure. A parallel structure for implementation of the encoder and decoder is possible for high data rate links.

The performance results of the coding and modulation schemes approach the theoretical limits over a Gaussian channel. This includes achievable performance even at high coding rates.

An important feature of the new scheme is its simplicity. The use of similar four-state encoder structures for the inner and the outer code leads to a significant reduction in implementation complexity. This is especially important in the case of high-speed decoders that require a highly parallelized architecture.

The frame structure governing the data transfer from the upper Data Link layers to the Physical layer is described. The data transfer frame structures at baseband and Physical layer are designed such that the data flow is compatible with the format used in the upper stream modules while providing support for maintaining synchronization during transition from one spectral efficiency mode to another. The reconfiguration of the modulation and coding is supported during operation without any interruption or significant delay in data transmission.

2 ARCHITECTURE OVERVIEW

A general view of the functional blocks of the proposed architecture for the synchronization and channel coding sublayer is presented in figure 2-1. The synchronization and channel coding sublayer provides the following functions:

- Transfer Frame data interfacing (buffering and ACM format adaptation);
- Transfer Frame synchronization and pseudo-randomization;
- Forward error coding based on serially concatenated convolutional coding;
- Physical Layer Framing, synchronization, and pseudo-randomization (optional).

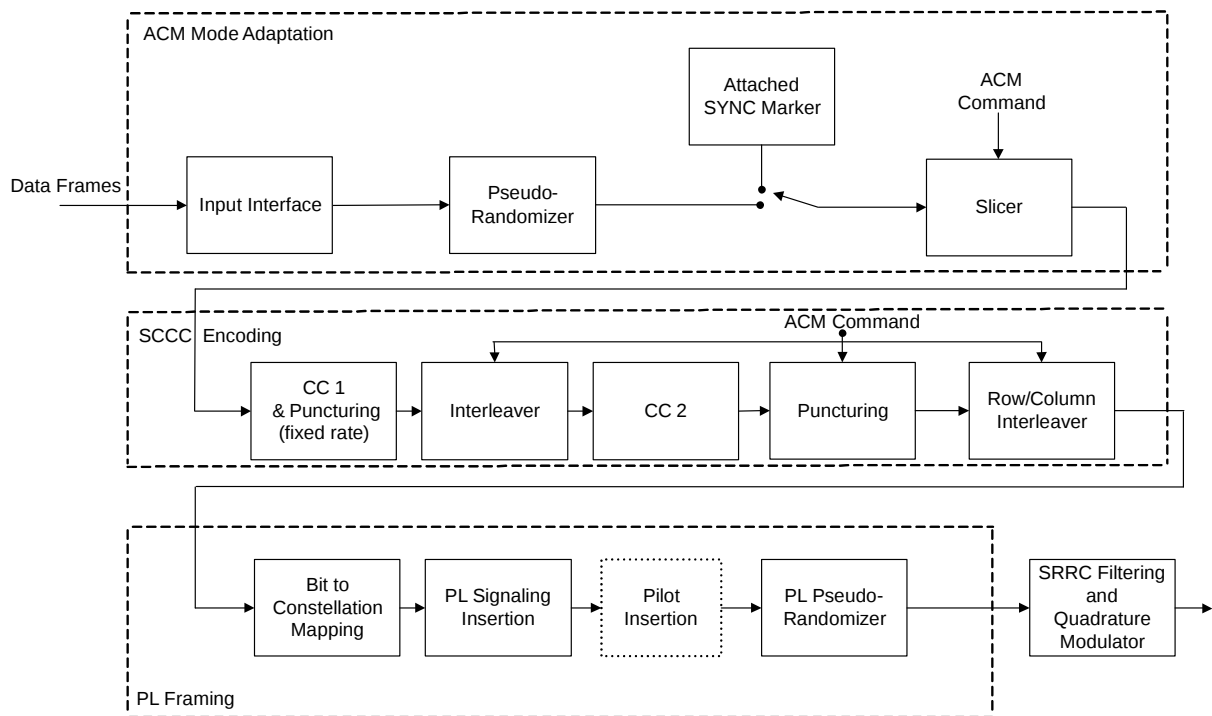


Figure 2-1: Functional Diagrams of Synchronization and Channel Coding Sublayer

Figure 2-2 illustrates the frame structures and stream formats at different stages of processing in synchronization and channel coding subsystems. The input data stream, transferred via the Transfer Frames, is compliant with proposed Data Link protocol in TM (reference [B4]) and AOS (reference [B5]).

Attached SYNC markers are inserted between Transfer Frames prior to encoding. The information blocks at the input of the encoder are formed by slicing the input data stream (after ASM insertion) into blocks of length K . The information block size varies depending on selected modulation and coding scheme (see table 4-3). A similar coding and modulation

scheme is applied to every 16 consecutive blocks that form a Physical Layer (PL) frame. The length of encoded blocks (N bits) is determined according to the modulation scheme (independent of the coding rate as shown in table 4-3). The length of encoded symbol blocks after encoding and mapping to modulation symbols is constant (8100 symbols) independent of the modulation and coding scheme. Maintaining constant symbol block sizes facilitates frame synchronization at the Physical layer.

Each subsystem is described in more detail in the following sections.

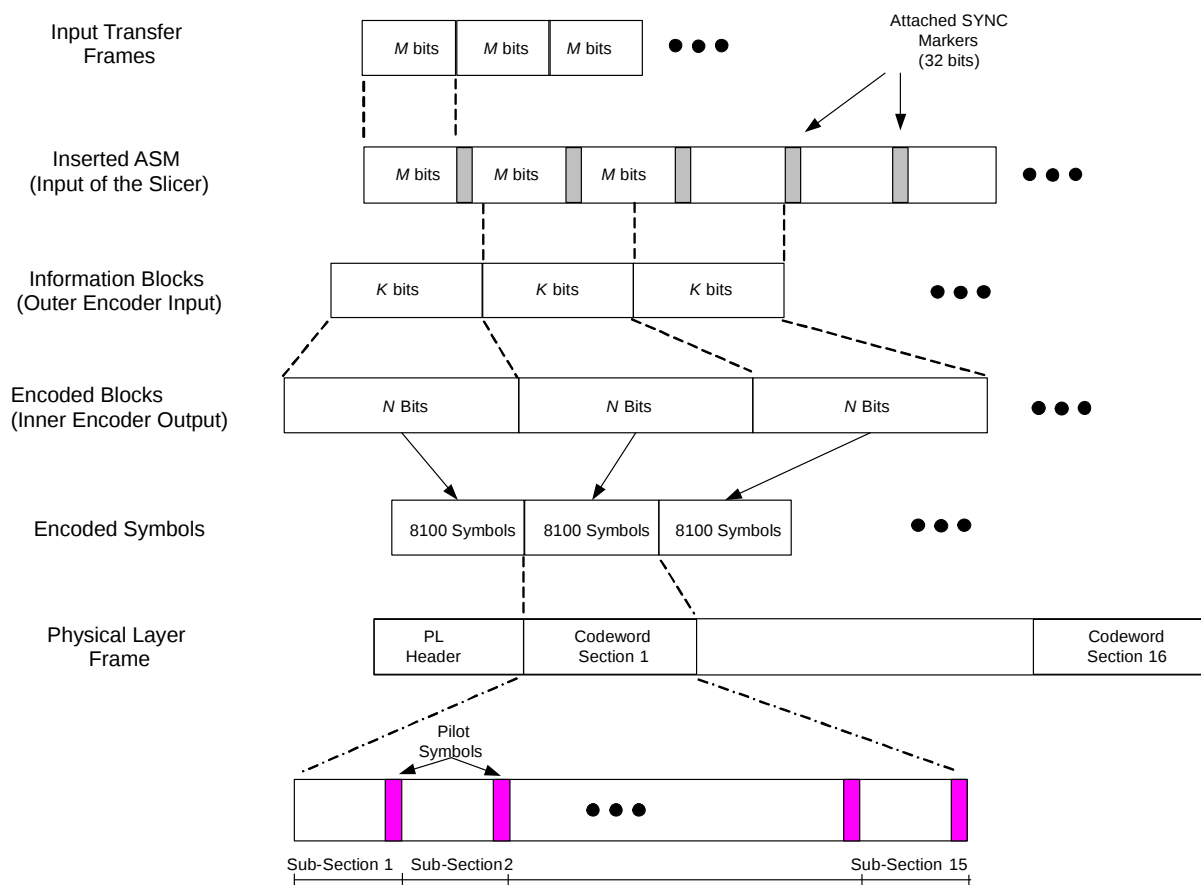


Figure 2-2: Stream Format at Different Stages of Synchronization and Channel Coding

3 MODE ADAPTATION

3.1 GENERAL

The mode adaptation unit provides the interface between the upstream units (data handling systems) and the on-board modulator. The input interface of this unit maps the input electrical format into a stream of logical bit format.

The data stream is passed through a pseudo-randomizer to ensure adequate bit transition density. In order to facilitate the frame synchronization at Data Link layer, a known sequence of bits, known as the Attached SYNC Marker (ASM), precedes each data Transfer Frame. The length of the ASM is fixed and the content is optimized to ensure a proper detection of the start of the Transfer Frame.

The Transfer Frame length at the input of the mode adaptation unit is predetermined. Neither the TM Space Data Link Protocol (reference [B4]) nor AOS Space Data Link Protocol (reference [B5]) specifies the Transfer Frame length. However, to maintain backward compatibility with the existing data link subsystems, the Transfer Frame length in bits (denoted as M in figure 2-2) is selected from one of the following values, as previously proposed in reference [B3]:

- 1784 bits (=223 x 1 octets);
- 3568 bits (=223 x 2 octets);
- 7136 bits (=223 x 3 octets);
- 8920 bits (=223 x 5 octets);
- 16384 bits.

The information block size is adjusted according to the spectral efficiency mode. A sub-module denoted as ‘slicer’ in figure 2-1 performs the mapping of the input bits (from the transfer data frame and ASM) into information blocks.

Mode Selection Mechanism

Several approaches for adjusting the modulation and coding schemes can be envisaged. One approach relies on the signal quality estimation (or prediction) provided by the ground station through a feedback channel (e.g., Telecommand). Alternative approach is to pre-schedule the change of modulation and coding schemes for each satellite pass based on geometrical information (elevation angle).

The choice of modulation and coding is communicated to the Synchronization and Coding modules via the ‘ACM Command’ signal. Using this command, the input data along with the SYNC markers are mapped into information blocks with a proper size to be processed by the outer encoder. Other functions such as interleaving, puncturing and bit-to-symbol mapping will also be adjusted according to the selected ACM format.

3.2 BASEBAND PSEUDO-RANDOMIZER

In order to ensure a proper detection of Transfer Frames at the receiver, the data stream must be sufficiently random. A pseudo-randomizer, similar to that described in reference [B3] is considered.

The pseudo-random sequence is applied starting with the first bit of the Transfer Frame. The Transfer Frame is randomized by exclusive-ORing the first bit of the Transfer Frame with the first bit of the pseudo-random sequence, followed by the second bit of the Transfer Frame with the second bit of the pseudo-random sequence, and so on. The pseudo-randomizer is not applied to the ASM. The attached SYNC marker is used to synchronise the pseudo-randomizer (reference [B3]). Figure 3-1 illustrates the pseudo-randomizer configuration.

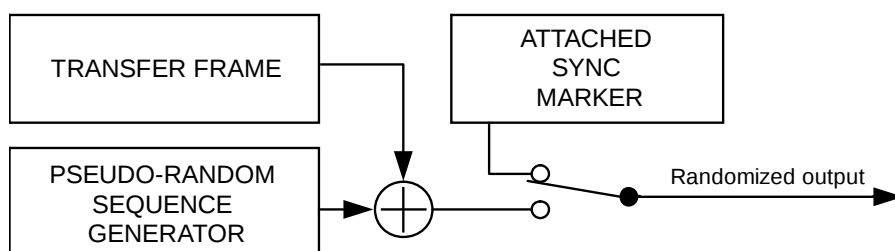


Figure 3-1: Pseudo-Randomizer Configuration

The pseudo-random sequence is generated using the following polynomial:

$$h(x) = x^8 + x^7 + x^5 + x^3 + 1 \quad (3-1)$$

This sequence begins at the first bit of the Transfer Frame and repeats after 255 bits, continuing repeatedly until the end of the Transfer Frame. The sequence generator is re-initialized to an all-ones state during each ASM period.

The first 40 bits of the pseudo-random sequence from the generator are shown below; the left-most bit is the first bit of the sequence to be exclusive-ORed with the first bit of the Transfer Frame; the second bit of the sequence is exclusive-ORed with the second bit of the Transfer Frame, and so on.

1111 1111 0100 1000 0000 1110 1100 0000 1001 1010

The pseudo-random sequence is added (Modulo 2) to the input Transfer Frame starting from the first bit of the Transfer Frame. Figure 3-2 shows the logical diagram of the pseudo-random sequence generator.

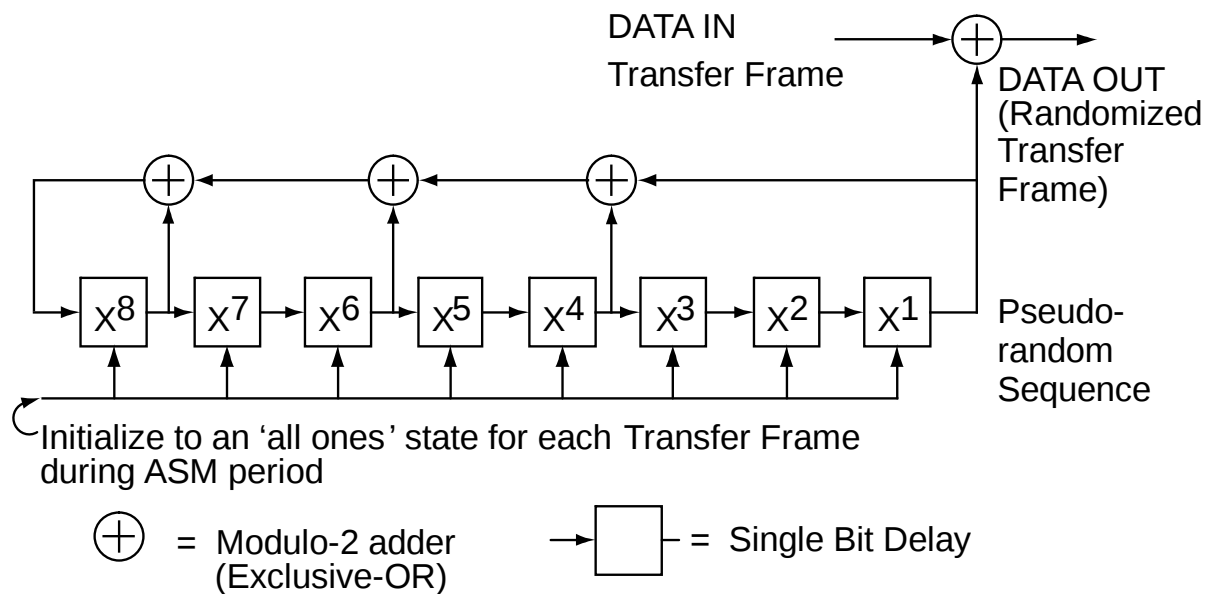


Figure 3-2: Pseudo-Randomizer Logic Diagram

3.3 ATTACHED SYNC MARKER

A 32-bit sequence of known bits is inserted between Transfer Frames in order to ensure frame synchronization at the receiver at the Data Link layer. The selected 32-bit ASM bit pattern in hexadecimal notation is shown in table 3-1 (similar to the 32-bit ASM pattern reported in reference [B3]).

Table 3-1: ASM Marker

ASM length	32 bits
ASM sequence (Hex)	1ACFFC1D

The same ASM length is used for all different Transfer Frames and all ACM formats. Since there are additional synchronization symbols at the Physical layer, the ASM sequence is mainly used to detect the start of the Transfer Frame. The ASM sequence is encoded along with the Transfer Frame. The probability of bit error in the ASM sequence at the nominal operating point is extremely low. Therefore, there is no need for a longer-than-32-bit ASM sequence for Data Link layer synchronization.

3.4 SLICER

The information block size at the input of SCCC encoder is determined based on the selected modulation and coding (ACM format). Prior to encoding of each information block, the input

Transfer Frames along with the attached SYNC markers are grouped into blocks of length K , corresponding to the information block size of the selected ACM format. No particular alignment between the Transfer Frame and the information blocks is considered. At the receiver side, the attached SYNC markers are used to reconstruct the Transfer Frames after decoding the information blocks.

The buffer size in the slicer should be dimensioned properly such that the switching between ACM formats is supported without insertion of a dummy frame or drop of any Transfer Frame. Table 3-2 summarizes the information block sizes of 27 selected ACM formats.

Table 3-2: Information Block Sizes for Different ACM Formats

ACM Format	Information Block Size (bits)	ACM Format	Information Block Size (bits)
1	5758	15	23518
2	6958	16	25918
3	8398	17	28318
4	9838	18	25918
5	11278	19	25318
6	13198	20	30958
7	11278	21	33358
8	13198	22	35998
9	14878	23	33358
10	17038	24	35998
11	19198	25	38638
12	21358	26	41038
13	19198	27	43678
14	21358		

4 SCCC ENCODER DESCRIPTION

4.1 OVERVIEW

Figure 4-1 shows the block diagram of the SCCC encoder. As described in 3.4, the input bits to the encoder are mapped into frames of size K . The frame size is selected according to the specific modulation order and coding rate (ACM format) with the objective of maintaining a constant length for modulated symbols per frame. The length of modulation symbols per frame is set to 8100 symbols.

Each frame is first encoded by an outer convolutional encoder at rate $1/2$ followed by a fixed puncturing scheme which sets the rate of the outer code to $r=2/3$. The resulting frame has $I=3/2 (K+2)$ encoded bits. The terminating bits are included in order to drive the code trellis to the zero state.

The intermediate frame of systematic and parity bits is interleaved according to an ad-hoc permutation law. An inner convolutional encoder has a coding structure that is similar to that of the outer code. The puncturing rule determines the actual SCCC code rate.

The SCCC scheme provides performance results close to the theoretical bounds. Such performance is achieved even at high coding rate. The good performance at high coding rate is due to the special arrangement of the encoding and the puncturing at the SCCC output. The impact of this structure at the receiver is that the information exchange between the two Soft-Input Soft-Output (SISO) modules in the iterative decoder corresponds to the cascade of a true rate $2/3$ encoder (the outer code) with a true rate $1/2$ encoder (the inner code). Subsequently, the interleaver gain that is typically larger for a lower outer code rates is maintained even in the case of a heavy puncturing of the outer code. This is mainly the reason for the good performance of the scheme, in terms of error floor, even at very high code rates.

In the following sections, each block of the SCCC encoder is described in more details.

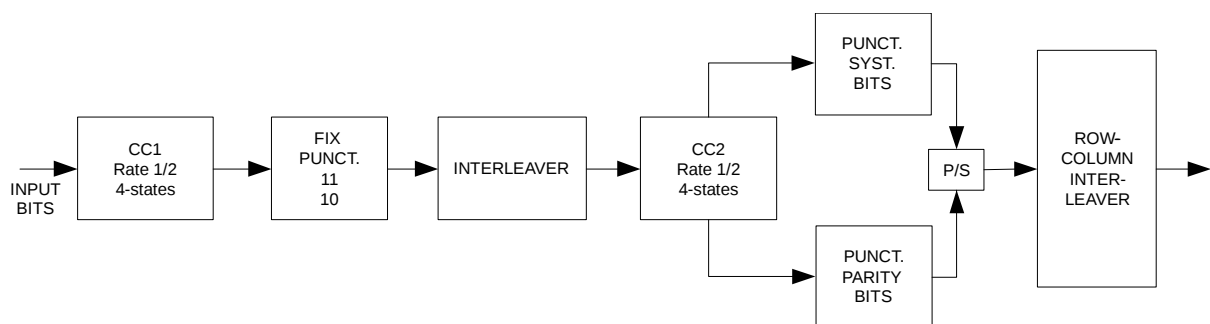


Figure 4-1: Block Diagram of the SCC Turbo Coding Scheme

4.2 CONVOLUTIONAL ENCODING

Both inner and outer convolutional encoders are based on the same four-state, rate 1/2 recursive, systematic encoder. The code structure is shown in figure 4-2.

For all SCCC rates the outer encoder is punctured to a rate 2/3 through the optimal puncturing scheme by maintaining all the systematic bits while decimating the parity bits by half, as shown in figure 4-3. The overall coding rate adjustment is carried out by puncturing the output bits of the inner code. A detailed description of the puncturing scheme is provided in 4.4.

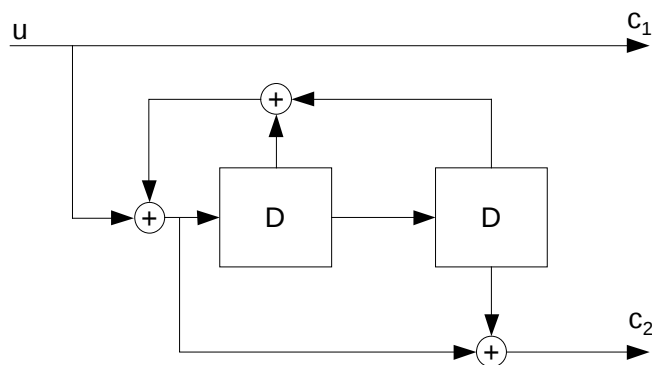


Figure 4-2: The Convolutional Encoder Block Diagram for CC1 and CC2

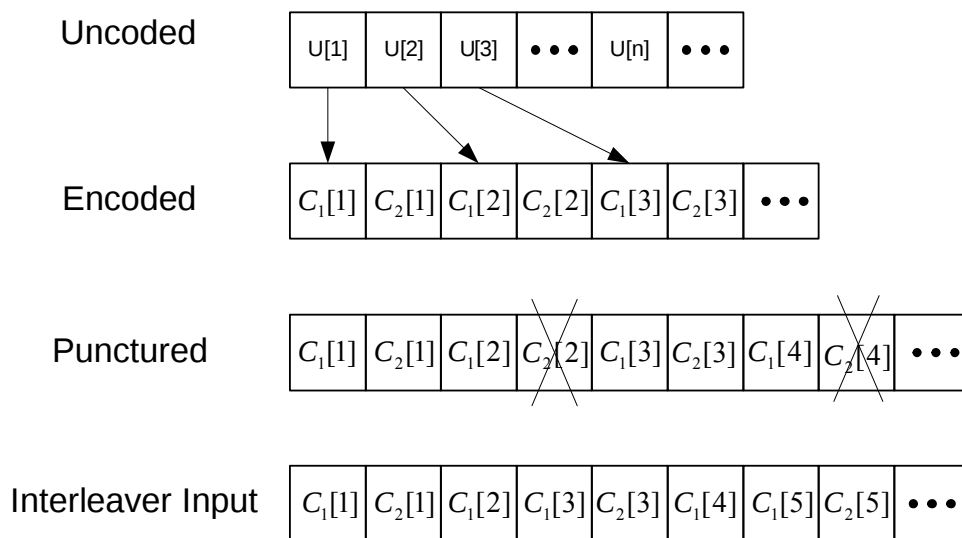


Figure 4-3: Outer Code Puncturing Scheme

4.3 INTERLEAVER

A proper selection of interleaver is a crucial part of the code design. The interleaver length I and the corresponding permutation law are selected according to the ACM format so as to keep the length of the frame after modulation mapping to a constant 8100 modulated symbols.

The interleaver is described by an ad-hoc permutation law which is reported in annex C. Table 4-1 lists interleaver sizes corresponding to different ACM formats. It is worth noting that for 27 selected ACM formats there are 19 different interleaver sizes.

Table 4-1: Interleaver Sizes for Different ACM Formats

ACM Format	Interleaver Length (bits)	ACM Format	Interleaver Length (bits)
1	8640	15	35280
2	10440	16	38880
3	12600	17	42480
4	14760	18	38880
5	16920	19	42480
6	19800	20	46440
7	16920	21	50040
8	19800	22	54000
9	22320	23	50040
10	25560	24	54000
11	28800	25	57960
12	32040	26	61560
13	28800	27	65520
14	32040		

4.4 CODING RATE ADJUSTMENT

4.4.1 GENERAL

In order to obtain the desired coding rate, puncturing is performed at the output of the inner encoder. The upper register at the output of the inner encoder (as shown in figure 4-1) contains the inner systematic bits, which coincide with the interleaved outer codeword as well as two additional bits terminating the inner trellis. The lower register, instead, contains the $I+2$ parity-check bits generated by the inner encoder. Two different puncturing algorithms are used to puncture bits in the upper and lower registers, corresponding to the inner systematic and parity bits accordingly.

Puncturing in the upper register is performed on the I inner systematic bits (excluding the two inner code terminating bits that are always transmitted) according to a periodic puncturing pattern. The puncturing period is based on 200 trellis steps, which corresponds to 300 outer coded bits.

It is worth noting that the systematic bits of the inner encoder correspond to the entire encoded bits generated by the outer encoder. The puncturing pattern of the systematic bits has been designed to maximise the free distance of the outer encoder, taking into account the impact of interleaving.

In practice, an exhaustive search for identifying optimal puncturing patterns is computationally infeasible. As a result, a suboptimal searching algorithm with manageable computation complexity has been devised. The penalty of using a sub-optimal puncturing pattern in the achievable performance is minimal. The search is performed incrementally, in a rate-compatible fashion, so that the punctured positions for a given rate are also punctured for all higher rates.

The optimisation of the upper register puncturing pattern involves both the number and the position of bits to be punctured. The outer code 'equivalent' rate, as seen from the SCCC output is approximated as:

$$R_o \approx \frac{200 \cdot n_p - 2}{S_{\text{sur}} \cdot n_p + 2} \quad (4-1)$$

where S_{sur} denotes the number of surviving bits in each 300-bit segment of the upper register after puncturing, and n_p is the number of puncturing segments in the interleaver, i.e., $n_p = I/300$.

Since n_p is not in general an integer number, the exact number S of transmitted systematic bits depends not only on the parameter S_{sur} but also on the interleaver size and on the puncturing pattern. The last two terminating bits of the inner encoder are always transmitted.

4.4.2 PUNCTURING SYSTEMATIC BITS

The design procedure for the outer code puncturing yielding outer code rates in the whole range $200/300 \leq R_o \leq 200/201$ can be described as follows:

- 1) Loop on all successive puncturing positions $i=1,\dots,99$ leading to a number of surviving systematic bits, out of 300, $S_{\text{sur}}=300-i$.
- 2) Loop on all S_{sur} remaining positions j that can be punctured in order to increase the outer code rate.
- 3) Set the new candidate position to j .
- 4) Compute the Weight-Enumerating Function (WEF) $A(d)$ of the outer encoder.
- 5) Save the position j and its WEF if it is better than the previous one. The position is 'better' if it has orderly small coefficients $A(1), A(2), \dots$ up to $A(8)$.
- 6) Consider the next candidate position.
- 7) Store the best position.
- 8) Next rate $200/S_{\text{sur}}$.

The search results for the optimal puncturing positions are reported in table 4-2. For each index i and parameter S_{sur} , the corresponding outer code rate (obtained by puncturing the bit in the position reported in the same row and all preceding ones) is given. From the implementation complexity point of view, the storage of the optimal puncturing pattern simply requires storing of the 100 puncturing positions reported in column 4 and 8 of table 4-2.

Table 4-2: Best Incremental Puncturing Positions

Index	S_{sur}	Rate	Punct. Pos.	Index	S_{sur}	Rate	Punct. Pos.
1	299	0.6689	76	51	249	0.8032	72
2	298	0.6711	1	52	248	0.8065	15
3	297	0.6734	145	53	247	0.8097	297
4	296	0.6757	214	54	246	0.8130	210
5	295	0.6780	256	55	245	0.8163	138
6	294	0.6803	37	56	244	0.8197	102
7	293	0.6826	109	57	243	0.8230	174
8	292	0.6849	181	58	242	0.8264	39
9	291	0.6873	277	59	241	0.8299	249
10	290	0.6897	235	60	240	0.8333	57
11	289	0.6920	55	61	239	0.8368	120
12	288	0.6944	127	62	238	0.8403	156
13	287	0.6969	163	63	237	0.8439	84
14	286	0.6993	19	64	236	0.8475	228
15	285	0.7018	199	65	235	0.8511	192
16	284	0.7042	91	66	234	0.8547	282
17	283	0.7067	289	67	233	0.8584	261
18	282	0.7092	244	68	232	0.8621	24
19	281	0.7117	64	69	231	0.8658	237
20	280	0.7143	268	70	230	0.8696	60
21	279	0.7168	223	71	229	0.8734	201
22	278	0.7194	136	72	228	0.8772	294
23	277	0.7220	172	73	227	0.8811	132
24	276	0.7246	28	74	226	0.8850	96
25	275	0.7273	100	75	225	0.8889	159
26	274	0.7299	190	76	224	0.8929	33
27	273	0.7326	10	77	223	0.8969	264
28	272	0.7353	46	78	222	0.9009	114
29	271	0.7380	118	79	221	0.9050	177
30	270	0.7407	154	80	220	0.9091	225
31	269	0.7435	81	81	219	0.9132	78
32	268	0.7463	207	82	218	0.9174	12
33	267	0.7491	259	83	217	0.9217	150
34	266	0.7519	292	84	216	0.9259	51
35	265	0.7547	232	85	215	0.9302	273
36	264	0.7576	67	86	214	0.9346	204
37	263	0.7605	280	87	213	0.9390	105
38	262	0.7634	247	88	212	0.9434	3
39	261	0.7663	147	89	211	0.9479	240
40	260	0.7692	30	90	210	0.9524	168
41	259	0.7722	111	91	209	0.9569	69
42	258	0.7752	183	92	208	0.9615	123
43	257	0.7782	6	93	207	0.9662	21
44	256	0.7813	48	94	206	0.9709	216
45	255	0.7843	93	95	205	0.9756	285
46	254	0.7874	165	96	204	0.9804	141
47	253	0.7905	129	97	203	0.9852	252
48	252	0.7937	219	98	202	0.9901	186
49	251	0.7968	195	99	201	0.9950	205
50	250	0.8000	270	100	200	1.0000	36

4.4.3 PUNCTURING PARITY BITS

4.4.3.1 General

Puncturing positions of the lower register, corresponding to the $I+2$ parity-check bits generated by the inner encoder, are determined using a rate-matching algorithm. This algorithm is described in detail in 4.4.3.2.

The puncturing of parity bits results in deleting a set of equally spaced bits. The number of deleted parity bits is determined based on the rate matching parameter Δ/I , representing the ratio between the number of deleted parity bits Δ and the overall number of parity bits I before puncturing:

$$\Delta = I - (P - 2) \quad (4-2)$$

where $P=N-S$ is the total number of transmitted parity check bits. It should be noted that the last two terminating parity check bits are always transmitted.

Some care must be taken in the choice of the parameter Δ , as explained below. The combination of the four-state trellis structure of the outer encoder and the fixed puncturing pattern of rate $2/3$ outer code results in a periodic encoder impulse response with a period of three bits.¹ If the parameter Δ is chosen such that the distance between two consecutive deleted parity bits is an integer multiple of three, the performance of the code will be degraded catastrophically. In fact, in this case there exist input sequences with weight 1 that will generate output sequences with weight 2. In other words, the inner encoder behaves as a feed-forward code, and the advantages of the recursive encoder structure disappear.

4.4.3.2 Rate Matching Algorithm

Given the two parameters Δ (number of bits to be deleted) and I (total number of bits), the rate-matching algorithm can be described using the following procedure:

- 1) Set the variable $e=1$.
- 2) For all possible positions i from 0 to $I-1$:
 - a) if $e>0$ transmit the i^{th} bit, else set $e=e+I$,
 - b) set $e=e-\Delta$;
- 3) Continue.

It should be noted that the last two terminating bits are always transmitted.

¹ The encoder impulse response is defined as the code sequence corresponding to a weight-1 input information sequence.

For each SCCC overall coding rate the optimisation of the coding scheme requires determining the parameter S_{sur} and the positions of the upper register punctured bits as reported in table 4-2. This first step is achieved via a mixed heuristic-simulation approach.

In each case, the value of S_{sur} determines the overall number of transmitted systematic bits S and, subsequently, the number of transmitted parity check bits P and the parameter Δ used by the rate-matching algorithm.

The optimisation campaign has been carried out for a set of desired SCCC rates. Table 4-3 shows required parameters that describe the encoder for 27 selected ACM formats.

Table 4-3: Main Encoder Parameters for 27 Selected ACM Formats

ACM format	m	S_{sur}	K	I	S	P	N	Δ
1	2	300	5758	8640	8642	7558	16200	1084
2	2	300	6958	10440	10442	5758	16200	4684
3	2	274	8398	12600	11510	4690	16200	7912
4	2	251	9838	14760	12351	3849	16200	10913
5	2	234	11278	16920	13200	3000	16200	13922
6	2	218	13198	19800	14390	1810	16200	17992
7	3	292	11278	16920	16470	7830	24300	9092
8	3	240	13198	19800	15842	8458	24300	11344
9	3	250	14878	22320	18602	5698	24300	16624
10	3	234	17038	25560	19939	4361	24300	21201
11	3	221	19198	28800	21218	3082	24300	25720
12	3	214	21358	32040	22857	1443	24300	30599
13	4	255	19198	28800	24482	7918	32400	20884
14	4	241	21358	32040	25741	6659	32400	25383
15	4	230	23518	35280	27051	5349	32400	29933
16	4	220	25918	38880	28515	3885	32400	34997
17	4	211	28318	42480	29880	2520	32400	39962
18	5	245	25918	38880	31755	8745	40500	30137
19	5	234	28318	42480	33137	7363	40500	35119
20	5	224	30958	46440	34677	5823	40500	40619
21	5	217	33358	50040	36197	4303	40500	45739
22	5	210	35998	54000	37802	2698	40500	51304
23	6	236	33358	50040	39366	9234	48600	40808
24	6	228	35998	54000	41042	7558	48600	46444
25	6	220	38638	57960	42507	6093	48600	51869
26	6	214	41038	61560	43915	4685	48600	56877
27	6	208	43678	65520	45429	3171	48600	62351

4.5 ROW-COLUMN INTERLEAVER

Prior to the bit-to-symbol mapping at the transmitter, a row-column interleaver is used to pseudo-randomize the selection of bits that are assigned to one modulation symbol. This is to ensure that the correlation between bits assigned to one symbol does not adversely affect the decoding process. This technique is known as Bit Interleaved Coded Modulation (BICM), as introduced in reference [B8].

Figure 4-4 illustrates the bit-interleaving scheme. The interleaver depth (number of rows) is equal to the size of one codeword section (i.e., 8100 symbols) and the number of columns is equal to m , where m is the modulation order.

The maximum memory size to implement the bit-interleaver is $m \times 8100 = 6 \times 8100 = 48600$ locations, each containing one bit, for the 64 APSK modulation scheme. The memory can be

seen as a matrix composed of m columns and 8100 rows. The number of rows is independent of the code rate and modulation scheme.

To implement the pragmatic code permutation, the output of the inner encoder, after puncturing, is bit interleaved using the block interleaving shown in figure 4-4. The input data is serially written into the interleaving column-wise and serially read out row-wise (the most significant bit is read out first).

Systematic bits (corresponding to the upper branch of the inner convolutional encoder) are first written sequentially in the register followed by the parity check bits (corresponding to the lower branch of the convolutional encoder). As a result of this scheduling, the systematic bits are mapped into the most significant bits and parity check bits to the least significant bits of the labels of the modulation symbols. The choice of modulation mapping, as described in 4.6, guarantees a higher protection for systematic bits.

The bit interleaving structure has 8100 rows, independent of the ACM format, and m columns where m is the modulation efficiency. The first symbol carries the bits positioned at index 0, 8100, 16200, 24300, 32400, 40500, 48600, for 64 APSK for instance. The second symbol carries bits at position 1, 8101, 16201, 24301, 32401, 40501, 48601 and so on up to the last symbol (carrying bits 8099, 16199, 24299, 32399, 40499, 48599).

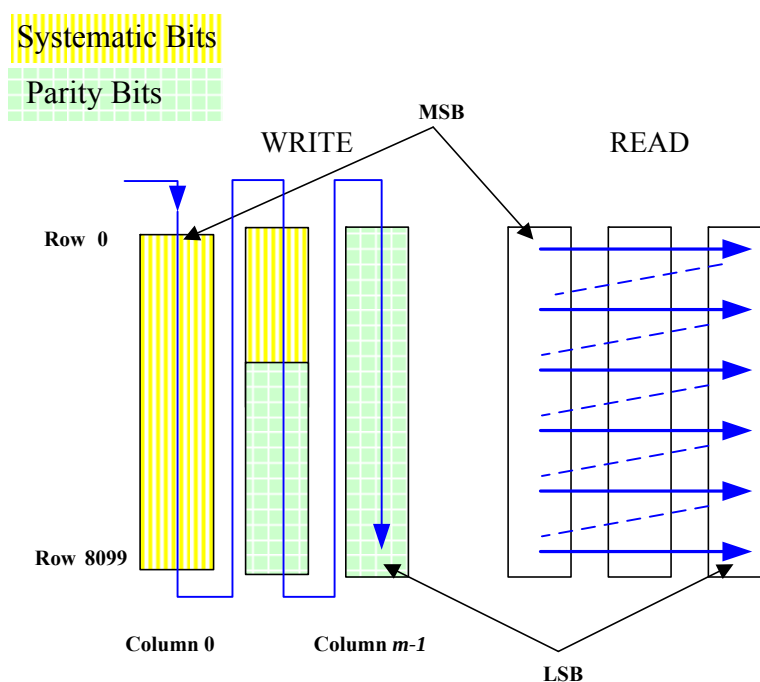


Figure 4-4: Row-Column Bit-Interleaving Scheme

4.6 BIT-TO-CONSTELLATION MAPPING

Figure 4-5 shows the selected modulation constellations along with the associated bits-to-symbols mapping laws (reference [B7]).

PSK Modulations

QPSK A conventional Gray-Coded QPSK modulation with absolute mapping (no differential coding) is considered. The normalized average energy per symbol is equal to one (Radius=1).

8PSK A conventional Gray-Coded 8PSK modulation with absolute mapping (no differential coding) is considered. The normalized average energy per symbol is equal to one (Radius=1).

APSK Modulations

The ratios between the radii of APSK constellations are denoted as $\gamma_1 = R_2/R_1$, $\gamma_2 = R_3/R_1$ and $\gamma_3 = R_4/R_1$. Figure 4-5 illustrates the bit-to-symbol mapping for APSK modulation schemes.

16APSK: This constellation is composed of two concentric circumferences, whose number of points is set to $N_1=4$ and $N_2=12$. Optimal values of γ_1 for 16APSK modulation schemes and linear channels are reported in table 4-4. Two possibilities for the normalization of modulated symbol are considered:

- the average signal energy is set to one: $[R_1]^2 + 3 [R_2]^2 = 4$;
- the radius of the outer constellation circle is set to one: $R_2=1$.

32APSK: This constellation is composed of three concentric circumferences whose number of points is set to $N_1=4$, $N_2=12$ and $N_3=16$. Optimal values of γ_1 and γ_2 for 32APSK modulation schemes are reported in table 4-4. Two possibilities for the normalization of modulated symbol are considered:

- the average signal energy is set to one: $[R_1]^2 + 3 [R_2]^2 + 4 [R_3]^2 = 8$;
- the radius of the outer constellation circle is set to one: $R_3=1$.

64APSK: This constellation is composed of four concentric circumferences, whose number of points is set to $N_1=4$, $N_2=12$, $N_3=20$ and $N_4=28$. For 64APSK modulation schemes, a set of parameters that maximize the minimum Euclidean distance are as follows: $\gamma_1=2.73$, $\gamma_2=4.52$ and $\gamma_3=6.31$. Further optimization of these parameters can be achieved by using similar methods reported in reference [B6] for linear and non-linear channels.

Two possibilities for the normalization of modulated symbol are considered:

- the average signal energy is set to one, $4[R_1]^2 + 12[R_2]^2 + 20[R_3]^2 + 28[R_4]^2 = 64$;
- the radius of the outer constellation circle is set to one: $R_4=1$.

Table 4-4: Optimum Constellation Radius Ratios for 16APSK and 32APSK

ACM Format	Modulation	Coding Rate	γ_1	γ_2
13	16-APSK	0.5925	3.82	N/A
14	16-APSK	0.6592	3.20	N/A
15	16-APSK	0.7259	2.95	N/A
16	16-APSK	0.7999	2.75	N/A
17	16-APSK	0.8740	2.60	N/A
18	32-APSK	0.6400	3.00	5.91
19	32-APSK	0.6992	2.95	5.65
20	32-APSK	0.7644	2.82	5.20
21	32-APSK	0.8237	2.65	4.67
22	32-APSK	0.8888	2.55	4.34

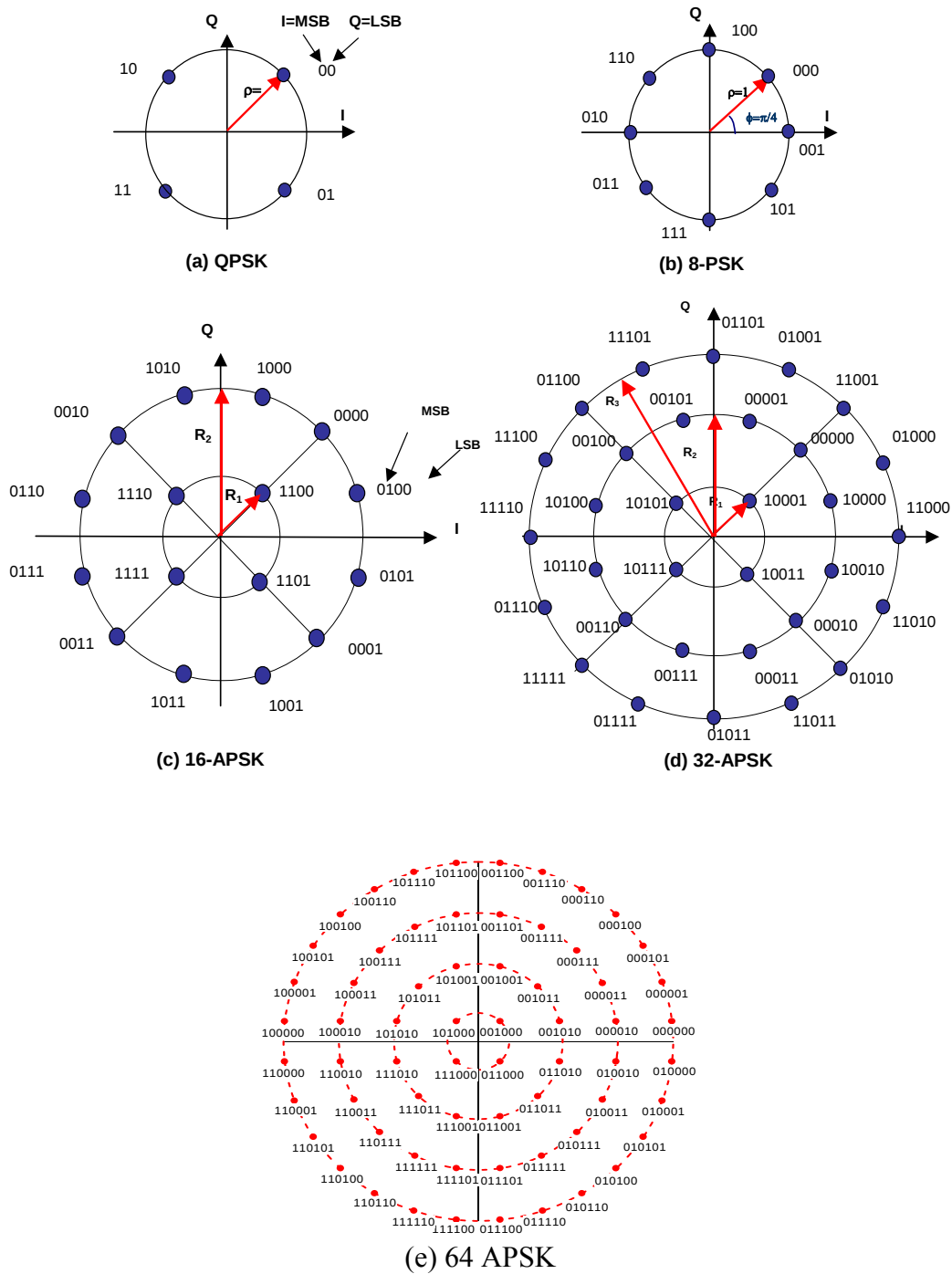


Figure 4-5: Bit Mapping into Constellations

4.7 SPECTRAL EFFICIENCY AND SUPPORTED SET OF ACM FORMATS

Table 4-5 presents a set of supported coding and modulation schemes (ACM formats) and related parameters. The rationale for selecting the set of ACM formats is highlighted below:

- The operating E_s/N_0 of selected formats covers steps of roughly one dB.
- The maximum coding rate is set to $R_c=9/10$.
- The maximum spectral efficiency is 5.4, corresponding to coding rate $r=9/10$ and 64-APSK modulation.
- The spectral efficiency from the unconstrained capacity formula for complex signals is given as:

$$\eta = \log_2 \left(1 + \frac{E_s}{N_0} \right) \quad (4-3)$$

Considering the maximum coding rate and equation 4-1, the modulation order can be determined as:

$$m = \left\lceil \frac{10}{9} \eta \right\rceil \quad (4-4)$$

- The number of *modulation symbols* in each codeword section is set to 8100 and imposed by the framing structure (this constraint simplifies frame synchronization significantly). This leads to a variable number of encoded bits N .
- A fixed size for the codeword section imposes constraints on the selection of nominal values of the SCCC encoder parameters (the information block size K , the interleaver length I , and the codeword size N).
- Constraints imposed on the interleaver length that depend on the architecture and degree of parallelism cause a slight change in the selected size of the interleaver I (to become I' as shown in table 4-5). Subsequently, the information block size K , the spectral efficiency and the theoretical E_s/N_0 (based on the capacity expression) change slightly as well. The resulting E_s/N_0 step is reported in the last column of table 4-5.

This approach defines the main 19 ACM formats in table 4-5, reported with regular font. The two highest spectral efficiencies for each modulation scheme have also been included with the modulation scheme with higher cardinality. This overlap is necessary since the coded-modulator performance can be different depending on the channel impairments.

In summary, a total of 27 ACM formats are supported, providing about 20 dB range in the required E_s/N_0 for the link budget. The expressions used in each column of table 4-5 are described below:

- eta: nominal spectral efficiency yielding the desired SNR according to the unconstrained capacity formula;

- K : nominal information block size;
- I : nominal interleaver length;
- N : number of encoded bits;
- m : modulation efficiency (defined as $\log_2$ of the modulator cardinality);
- E_s/N_0 : nominal signal-to-noise ratio thresholds derived from unconstrained Shannon Capacity Bound and nominal spectral efficiency η (η):

$$\frac{E_s}{N_0} = 10 \log_{10}(2^\eta - 1) \quad (4-5)$$

- *code rate*: overall rate of the code (K/N);
- K' : Information block size obtained applying the constraints on the interleaver block size that derives from parallelism.
- I' : Interleaver size applying the constraints that derive from required parallelism.
- η' : actual spectral efficiency, applying the constraints on the interleaver block size that derive from parallelism.
- E_s/N_0' : signal-to-noise ratio thresholds derived from the capacity bound by applying the constraints on the interleaver block size due to parallelism.

$$\left(\frac{E_s}{N_0}\right)' = 10 \log_{10}(2^{\eta'} - 1) \quad (4-6)$$

- *delta*: actual signal-to-noise ratio step.

Table 4-5: ACM Formats of the SCCC Encoder

Nominal values									Interleaver constrained				
	Es/N0	eta	K	I	N	m	Code Rate		K'	I'	eta'	Es/N0'	delta
QPSK	1	-1.85	0.7254	5,876	8,816	16,200	2	0.36	5,758	8,640	0.7109	-1.96	
	2	-0.85	0.8659	7,014	10,523	16,200	2	0.43	6,958	10,440	0.8590	-0.89	1.07
	3	0.15	1.0254	8,306	12,461	16,200	2	0.51	8,398	12,600	1.0368	0.22	1.11
	4	1.15	1.2039	9,752	14,630	16,200	2	0.60	9,838	14,760	1.2146	1.21	0.99
	5	2.15	1.4012	11,350	17,027	16,200	2	0.70	11,278	16,920	1.3923	2.11	0.90
	6	3.15	1.6164	13,092	19,640	16,200	2	0.81	13,198	19,800	1.6294	3.21	1.10
8PSK	7	2.15	1.4012	11,350	17,027	24,300	3	0.47	11,278	16,920	1.3923	2.11	-1.10
	8	3.15	1.6164	13,092	19,640	24,300	3	0.54	13,198	19,800	1.6294	3.21	1.10
	9	4.15	1.8484	14,972	22,460	24,300	3	0.62	14,878	22,320	1.8368	4.10	0.89
	10	5.15	2.0958	16,976	25,466	24,300	3	0.70	17,038	25,560	2.1035	5.18	1.08
	11	6.15	2.3568	19,090	28,637	24,300	3	0.79	19,198	28,800	2.3701	6.20	1.02
	12	7.15	2.6299	21,302	31,955	24,300	3	0.88	21,358	32,040	2.6368	7.18	0.98
16APSK	13	6.15	2.3568	19,090	28,637	32,400	4	0.59	19,198	28,800	2.3701	6.20	-0.98
	14	7.15	2.6299	21,302	31,955	32,400	4	0.66	21,358	32,040	2.6368	7.18	0.98
	15	8.15	2.9133	23,598	35,399	32,400	4	0.73	23,518	35,280	2.9035	8.12	0.94
	16	9.15	3.2056	25,966	38,951	32,400	4	0.80	25,918	38,880	3.1998	9.13	1.01
	17	10.15	3.5053	28,392	42,590	32,400	4	0.88	28,318	42,480	3.4960	10.12	0.99
32APSK	18	9.15	3.2056	25,966	38,951	40,500	5	0.64	25,918	38,880	3.1998	9.13	-0.99
	19	10.15	3.5053	28,392	42,590	40,500	5	0.70	28,318	42,480	3.4960	10.12	0.99
	20	11.15	3.8111	30,870	46,307	40,500	5	0.76	30,958	46,440	3.8220	11.19	1.07
	21	12.15	4.1220	33,388	50,084	40,500	5	0.82	33,358	50,040	4.1183	12.14	0.95
	22	13.15	4.4370	35,940	53,912	40,500	5	0.89	35,998	54,000	4.4442	13.17	1.03
64APSK	23	12.15	4.1220	33,388	50,084	48,600	6	0.69	33,358	50,040	4.1183	12.14	-1.03
	24	13.15	4.4370	35,940	53,912	48,600	6	0.74	35,998	54,000	4.4442	13.17	1.03
	25	14.15	4.7555	38,520	57,782	48,600	6	0.79	38,638	57,960	4.7701	14.20	1.02
	26	15.15	5.0766	41,120	61,682	48,600	6	0.85	41,038	61,560	5.0664	15.12	0.92
	27	16.15	5.4000	43,740	65,612	48,600	6	0.90	43,678	65,520	5.3923	16.13	1.01

5 PHYSICAL LAYER FRAMING

Figure 5-1 illustrates the Physical layer frame structure that consists of the following segments:

- Frame Header which consists of two fields:
 - Frame Marker: 256 known symbols used for start of frame detection and synchronization;
 - Frame Descriptor: 64 symbols carrying 7 bits of information to identify the ACM format used per each physical frame as well as the presence or absence of pilot symbols;
- Codeword Sections: 16 sections of modulation symbols (with additional pilot symbols as option).

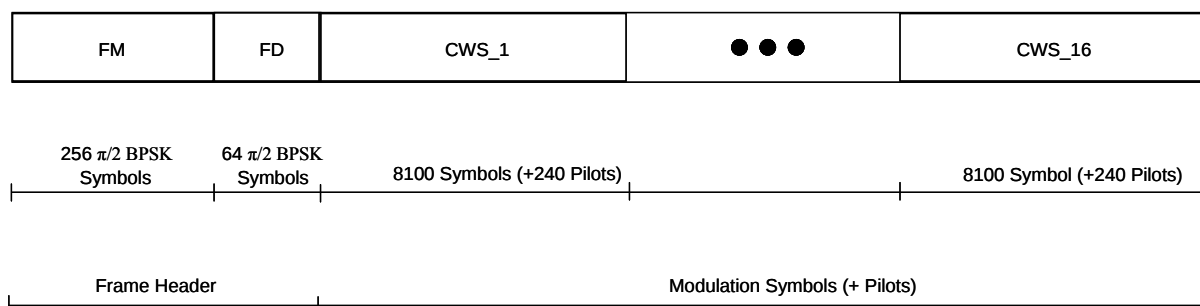


Figure 5-1: Physical Layer Frame Structure

5.1 FRAME MARKER

The Frame Marker consists of 256 known bits mapped into 256 $\pi/2$ -BPSK modulated symbols. The Frame Marker is used for detecting the start of the frame as well as initial timing and coarse carrier synchronization. The length and the modulating bit sequence of the Frame Marker are selected such that the start of frame can be detected with a low probability of detection error (misdetection as well as false alarm) in the presence of severe channel impairments.

The bit pattern used in the Frame Marker is generated using a Gold sequence. Figure 5-2 shows the logical block diagram of the proposed sequence generator using shift registers and exclusive-OR operators. The generator polynomials used for the feedback loop are:

$$\begin{aligned}
 g_1(x) &= x^8 + x^6 + x^5 + x^4 + 1 \\
 g_2(x) &= x^8 + x^6 + x^5 + x^4 + x^3 + x + 1
 \end{aligned}
 \tag{5-1}$$

The upper and the lower shift registers are initialized as shown in figure 5-2. The first 40 bits of the Frame Marker sequence for the generator are shown below. The left-most bit corresponds to the first modulating bit of the Frame Marker:

1111 1011 0100 0100 0001 1111 0001 1101 1011 1101 ...

The mapping of the Frame Marker sequence into $\pi/2$ -BPSK symbols is described in the next section.

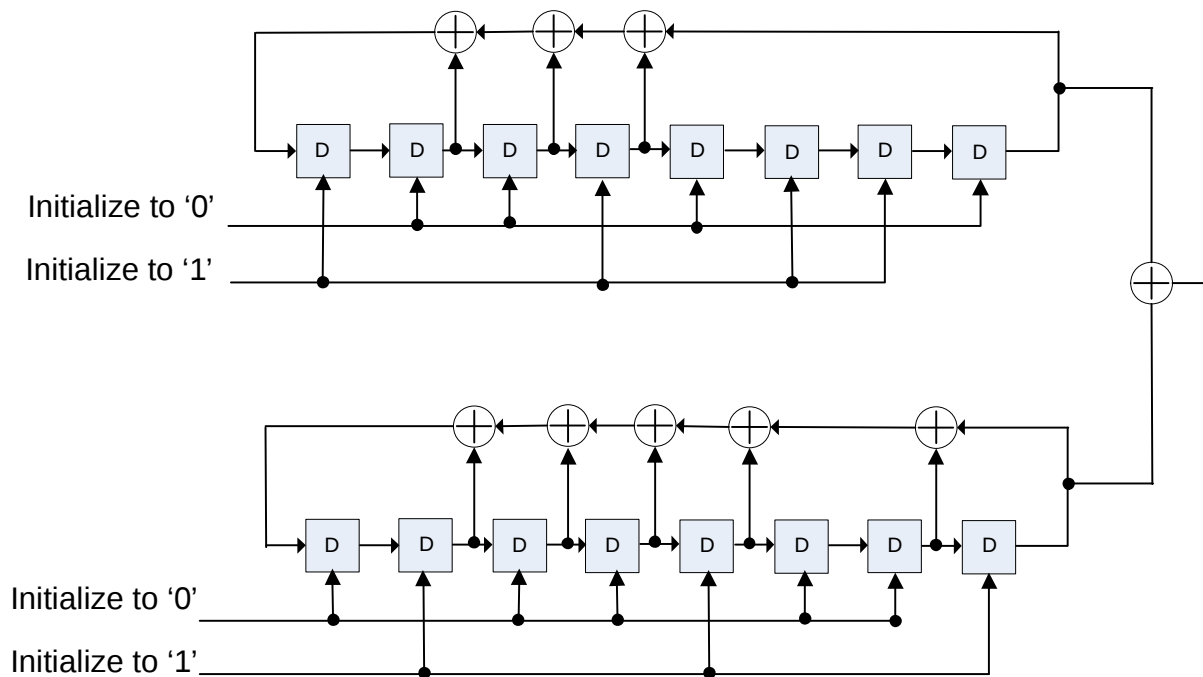


Figure 5-2: Frame Marker Sequence Generator

5.2 FRAME DESCRIPTOR STRUCTURE

5.2.1 GENERAL

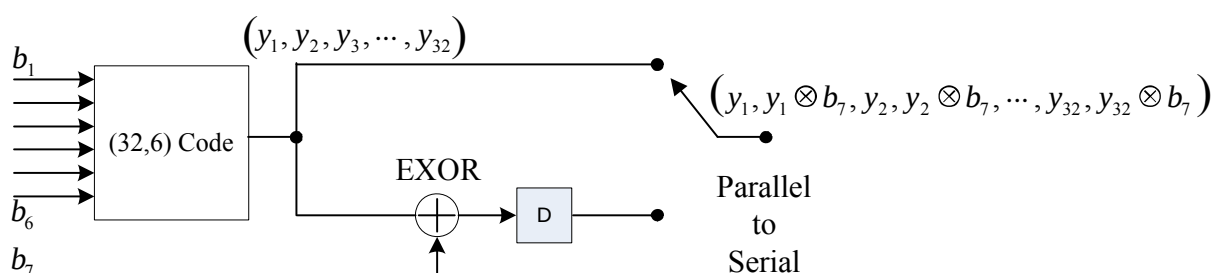
The Frame Descriptor carries seven bits of information. These bits identify the ACM format of codeword sections within a PL frame (five bits) as well as the absence or presence of distrusted pilots.

Information bits in the Frame Descriptor are encoded using a non-symmetric binary code of length 64 and dimension 7 with minimum distance $d_{\min}=32$. The code is similar to that used in reference [B2] for Physical layer signalling. The content of the Frame Descriptor is shown in table 5-1.

Table 5-1: Frame Descriptor Information Content

Bit Number	Content
b_1 - b_5	ACM Formats (Decimal values 1 to 27 are used)
b_6	Distributed Pilot On (=1) / Off (=0)
b_7	Reserved (Default =0)

The code is constructed using a bi-orthogonal (32,6) code and extended as shown in figure 5-3.

**Figure 5-3: Frame Descriptor Code Structure**

The first six bits, b_1 - b_6 , are encoded using a linear block code of length 32 with the following generator matrix:

$$G = \begin{bmatrix} 010101010101010101010101010101 \\ 00110011001100110011001100110011 \\ 00001111000011110000111100001111 \\ 00000000111111110000000011111111 \\ 00000000000000001111111111111111 \\ 11111111111111111111111111111111 \end{bmatrix}$$

Figure 5-4: Generator Matrix for (32,6) Code

The most significant bit is multiplied with the first row of the matrix, the following bit with the second row, and so on. The 32 coded bits is denoted as $(y_1 y_2 \dots y_{32})$. By default, the least significant bit of the Frame Descriptor is 0 and the final output is $(y_1 y_1 y_2 y_2 \dots y_{32} y_{32})$; i.e., each symbol shall be repeated. The 64-bit output code is further scrambled by the following binary sequence:

011100011001110110000011110010010101001101000010001011011111010.

5.2.2 FRAME HEADER MODULATION

The frame header, which consists of the Frame Marker (256 bits) and Frame Descriptor (64 bits), is modulated into 320 $\pi/2$ -BPSK symbols with the following mapping (similar to that of PL Header in reference [B2]).

Assuming that the Frame Header binary sequence is denoted as $(x_1 x_2 \dots x_{320})$, the In-phase (I) and Quadrature (Q) components of 320 $\pi/2$ -BPSK modulated symbols are determined according the following rule:

$$\begin{aligned} I_{2i-1} &= Q_{2i-1} = \frac{1}{\sqrt{2}}(1 - 2x_{2i-1}) \\ I_{2i} &= -Q_{2i} = \frac{-1}{\sqrt{2}}(1 - 2x_{2i}) \end{aligned} \quad \text{for } i=1, 2, \dots, 320 \quad (5-2)$$

5.3 PILOTS INSERTION

The use of distributed pilot symbols in codeword sections is considered to be an option to facilitate carrier and phase synchronization.

Figure 5-5 illustrates the proposed format for insertion of distributed pilots. The presence or absence of pilot symbols can be changed on a frame-by-frame basis using one bit (b_6) in the Frame Descriptor field (see table 5-1). Each pilot is an unmodulated symbol, with equal In-phase and Quadrature components: $I=(1/\sqrt{2})$, $Q=(1/\sqrt{2})$.

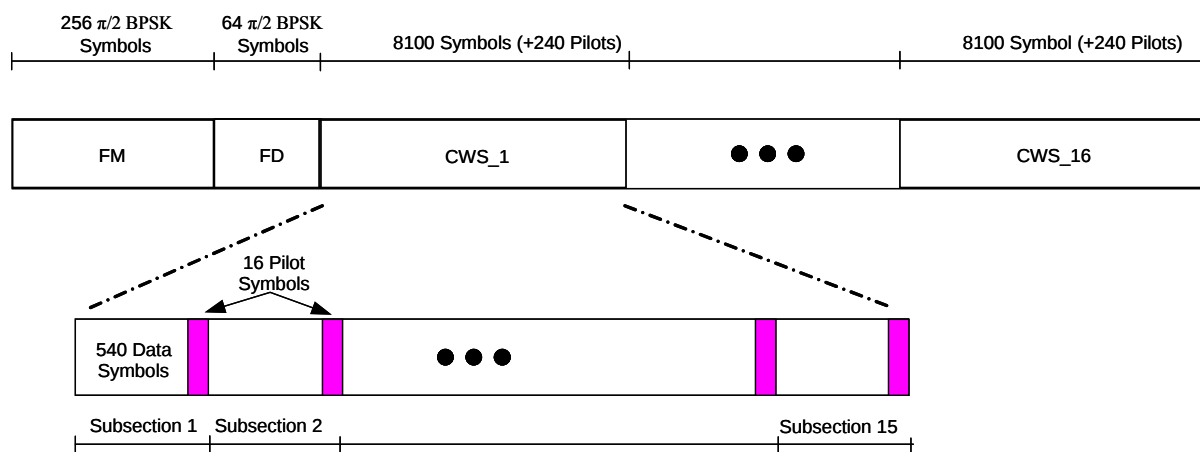


Figure 5-5: Distributed Pilot Pattern

Required parameters to specify the pilot distribution pattern within each PL frame are presented in table 5-2. The total overhead due to pilot insertion is around three percent.

Table 5-2: Frame Parameters Related to Pilot Distribution

Parameter	Value
Codeword Section Length	8100 symbols
Number of Codeword Section per frame	16 Sections
Number of sub-sections	15 sub-sections
Number of Data Symbols per sub-section	540 Symbols
Number pilots per subsection	16 Symbols
Total Number of pilots per section	240 Symbols
Total Section length including pilot symbols	8340 Symbols

5.4 PHYSICAL LAYER I/Q PSEUDO-RANDOMIZATION (OPTIONAL)

Physical layer randomization of codeword sections, including the data symbols as well as the pilots, is considered to be an option but not configurable during operation. The Physical layer randomization is used to disperse the signal energy in order to avoid any spectral spur due to repetitive data or pilot patterns. The proposed Physical layer pseudo-randomizer is similar to that in DVB-S2 standard (reference [B2]). For completeness, the PL pseudo-randomizer is described in annex C.

5.5 BASEBAND FILTERING

The baseband pulse shaping filter applied to In-phase and Quadrature signals is a square-root raised cosine filter:

$$H(f) = \begin{cases} 1 & |f| < f_N(1-\alpha) \\ \left\{ \frac{1}{2} + \frac{1}{2} \sin \frac{\pi}{2f_N} \left(\frac{f_N - |f|}{\alpha} \right) \right\}^{1/2} & f_N(1-\alpha) < |f| < f_N(1+\alpha) \\ 0 & |f| > f_N(1+\alpha) \end{cases} \quad (5-3)$$

where $f_N = \frac{1}{2T_s} = \frac{R_s}{2}$ is the Nyquist frequency and α is the roll-off factor.

The roll-off factor is fixed per operation and is selected from one of the following values: $\alpha=0.2, 0.25, 0.30$, or 0.35 .

6 ENCODER IMPLEMENTATION FOR HIGH BIT RATE DATA TRANSFER

In order to achieve a net data rate as high as 1 Gbits/sec, the encoder architecture (as well as the decoder which will not be described here) must be designed with some degree of parallelism. The trivial solution based on using n duplicated encoders (as shown in figure 6-1), each working at a reduced speed of R_S/n , has the drawback of requiring a memory n times larger than that of a single encoder. This amount of extra memory that is required for n parallel operations is significantly large in the case of large codeword sizes.

There is no obvious scheme to generate a parallel encoder without a full functional replication (e.g., acting on a single working memory). In fact, due to the inherent sequential nature of the encoding process, if the information block is split into n segments, the encoding of each segment will depend on the final state reached by the trellis in the previous one. However, there exists a solution by exploiting the properties of a linear system and encoding twice the information bits contained in each segment.

The first encoding procedure is used to evaluate the final states of each segment in which an information block is split, starting each encoder from the zero state. Next, by using the mathematical computation of the final states exploiting the linear properties of the encoder, the same data is re-encode with correct initial states obtained by a simple calculation.

The proposed scheme requires a modest increase in memory requirement, but it needs $2n$ encoder processors for n degree of parallelism. The additional cost of the encoders is quite marginal (two registers per encoder plus a limited combinatorial logic). This solution does not impose any constraint on the decoder.

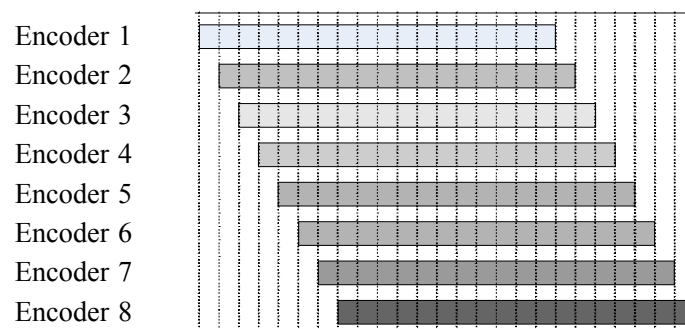


Figure 6-1: Encoder Functional Partitioning

The convolutional encoder described in 4.2 can be expressed using state equations:

$$\begin{cases} x(n+1) = A \bullet x(n) + B \bullet u(n) \\ y(n) = C \bullet x(n) + D \bullet u(n) \end{cases} \quad (6-1)$$

where the state, output, and input are respectively denoted as:

$$x(n) = \begin{bmatrix} x_0 \\ x_1 \end{bmatrix}; \quad y(n) = \begin{bmatrix} c_0 \\ c_1 \end{bmatrix}; \quad u(n) = u; \quad (6-2)$$

The dots ($\bullet$) represent matrix product defined over a Galois Field GF(2). Equivalently, the following modulo-2 arithmetic with the corresponding sum and product operators are defined:

$$\begin{cases} a, b \in (0,1) \\ a + b = \text{mod}(a + b, 2) \\ a * b = a \cdot b \end{cases} \quad (6-3)$$

in which $\text{mod}(\cdot)$ represents the modulo reduction operator.

The following matrixes describe the particular encoder implementation, in which A is the state transition matrix (with zero input):

$$A = \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}; B = \begin{bmatrix} 0 \\ 1 \end{bmatrix}; C = \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix}; D = \begin{bmatrix} 1 \\ 1 \end{bmatrix}; \quad (6-4)$$

Equation 6-1 can also be written recursively:

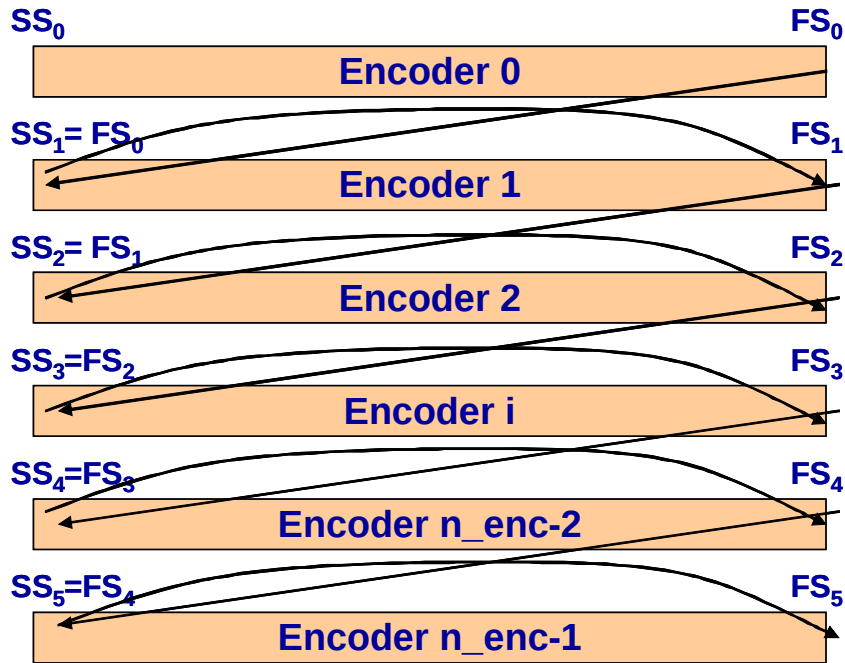
$$\begin{cases} x(n+1) = A \bullet x(n) + B \bullet u(n) = \\ = A^{n+1} \bullet x(0) + A^n B \bullet u(0) + A^{n-1} B \bullet u(1) + \dots + A^2 B \bullet u(n-2) + AB \bullet u(n-1) + B \bullet u(n) \\ y(n) = C \bullet (A \bullet x(n-1) + B \bullet u(n-1)) + D \bullet u(n) = \\ = C \bullet A^n \bullet x(0) + CA^{n-1} B \bullet u(0) + \dots + CAB \bullet u(n-2) + CB \bullet u(n-1) + D \bullet u(n) \end{cases} \quad (6-5)$$

In equation 6-5 one can identify a component of the state that depends only on previous states and another component that depends only on the input. The same applies to the output. Starting from the zero state, equation 6-5 can be turned to the following set of equations, defining the intermediate initial starting states (see also figure 6-2):

$$\begin{aligned}
x_0(0) &= \begin{bmatrix} 0 \\ 0 \end{bmatrix}; \\
x_1(0) &= FS_0; \\
x_2(0) &= FS_1^* = A^n \cdot FS_0 + FS_1; \\
x_3(0) &= FS_2^* = A^n \cdot FS_1^* + FS_2 = A^{2n} \cdot FS_0 + A^n \cdot FS_1 + FS_2; \\
x_4(0) &= FS_3^* = A^n \cdot FS_2^* + FS_3 = A^{3n} \cdot FS_0 + A^{2n} \cdot FS_1 + A^n \cdot FS_2 + FS_3; \\
x_5(0) &= FS_4^* = A^n \cdot FS_3^* + FS_4 = A^{4n} \cdot FS_0 + A^{3n} \cdot FS_1 + A^{2n} \cdot FS_2 + A^n \cdot FS_3 + FS_4;
\end{aligned} \tag{6-6}$$

Equation 6-6 constitutes the rule underneath the suggested encoder parallel operation, obtained with no memory expansion, compared with the traditional single encoder structure. The resulting architecture is shown in figure 6-3.

In equation 6-6 FS denotes the Final State and the indexes refer to the codeword partitioning shown in figure 6-2. Hence, a component of the final states, the one depending on input data only (and zero input state), is obtained by the pre-encoding process (first encoder operation); the second component of the final state is obtained by matrix exponentiation. The latter turns out to be very simple, due to the particular properties of the matrix A (see equation 6-7). Next, the final states are used as starting states of the following blocks (see diagonal arrows).



The Starting State and the Final State of encoder i is indicated with SS_i and FS_i , respectively. The total number of encoders is n_{enc} .

Figure 6-2: State Propagation Diagram

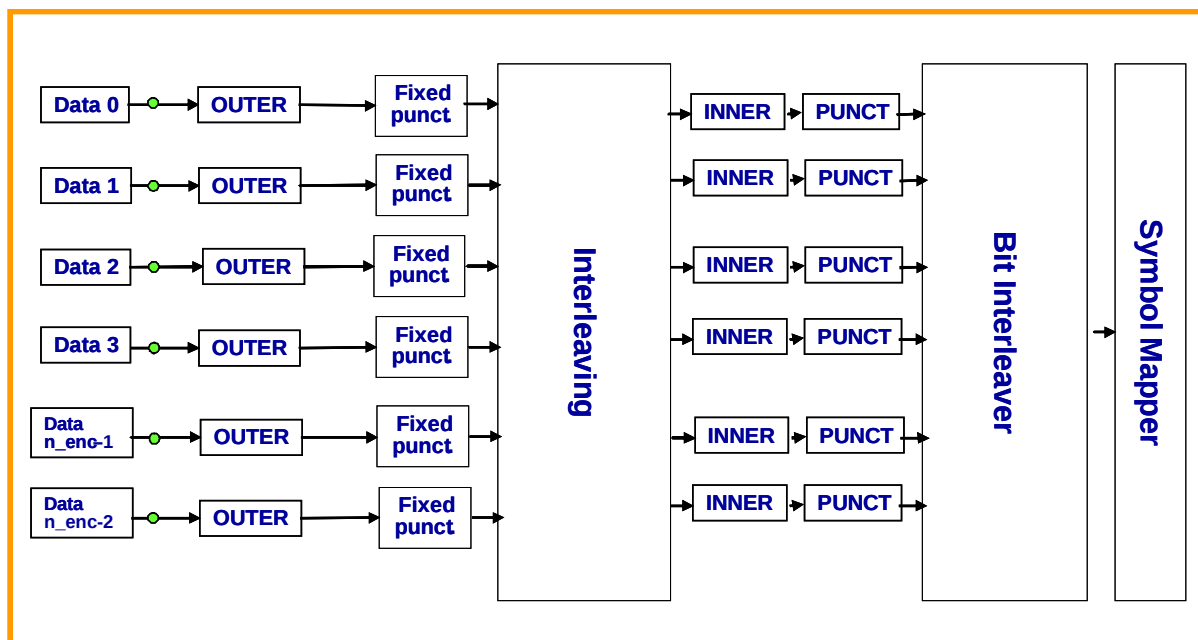


Figure 6-3: Parallelized Encoder Structure

The cyclic properties of the matrix A that can be exploited in the calculations are:

$$\begin{aligned}
 A &= \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}; \\
 A^2 &= \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix} \bullet \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}; \\
 A^3 &= A^2 \bullet A = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \bullet \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} = I; \\
 A^4 &= A^3 \bullet A = I \bullet A = A;
 \end{aligned} \tag{6-7}$$

To obtain a parallel SCCC encoder, we need a parallelism order equal to twice that normally necessary since we have to encode the same data twice. The maximum data rate at the inner encoder is 1.5 Gb/sec; hence, considering as an example the maximum speed of the encoder processor (in trellis steps/sec) equal to, $F_{ck} = 250$ MHz the number of parallel encoders, n_{enc} , can be calculated as:

$$n_{enc} = \frac{3 \text{ Gbit}}{250 \text{ MHz}} = 12;$$

Alternatively, if the clock frequency is set to $F_{ck} = 250$ MHz, the number of parallel encoders.

$$n_{enc} = \frac{3 \text{ Gbit}}{200 \text{ MHz}} = 15.$$

7 PERFORMANCE RESULTS

The performance results measured in terms of Bit Error Ratio (BER) and Frame Error Ratio (FER) for 27 selected combinations of coding and modulation schemes in AWGN channel are presented in figures 7-1 and 7-2, respectively.

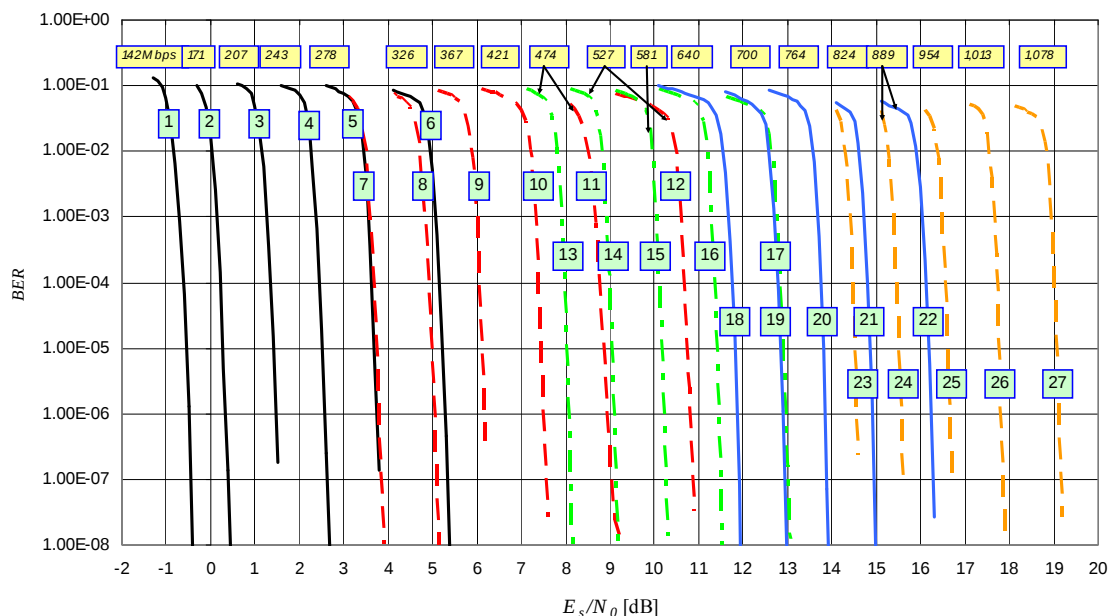


Figure 7-1: Simulated BER Performance of the 27 ACM Formats

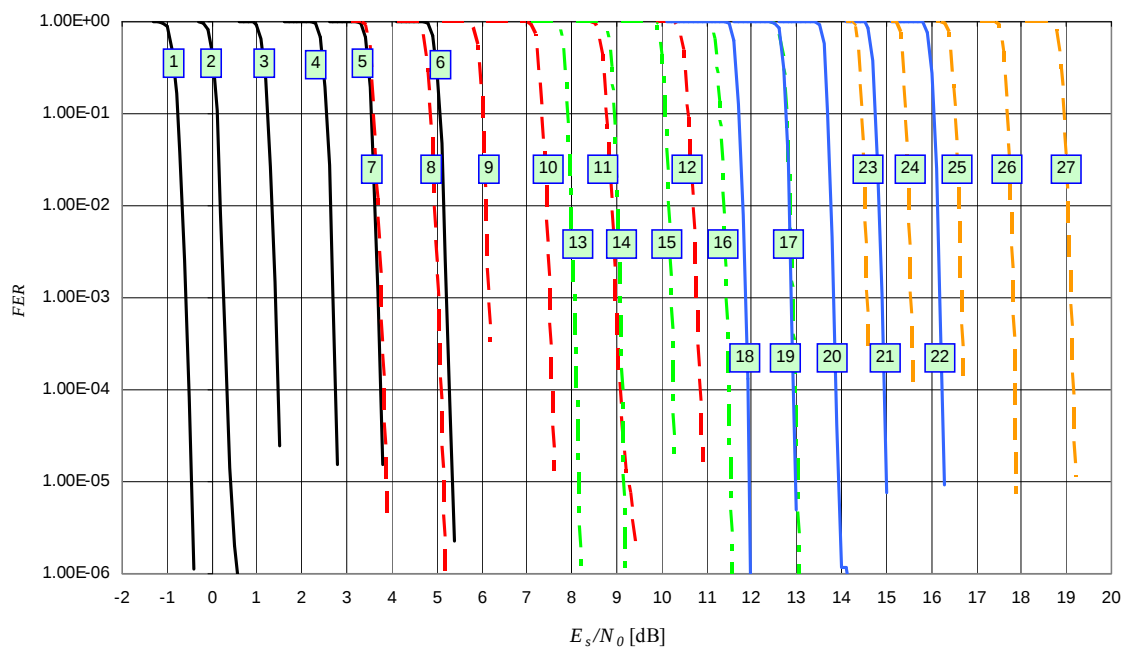


Figure 7-2: Simulated Frame Error Ratio of the 27 ACM Formats

The coding rates corresponding to 27 ACM formats can be determined based on the information block size K and the encoded block sizes N , as specified in table 4-3. As an example, for the ACM format No. 6 with $(K,N)=(13198,16200)$ and QPSK modulation, the required E_s/N_0 to achieve a BER of 10^{-8} is 5.4 dB. The corresponding Frame Error Ratio in this case is around 10^{-6} .

It should be noted that the computed Frame Error Ratio in figure 7-2 refers to the probability of occurrence of at least one erroneous bit in a decoded information block. The structure of the information block is illustrated in figure 2-2. As shown, there is no one-to-one mapping between the information blocks and the input Transfer Frames.

The information data rate labelling the curves in figure 7-1 corresponds to a symbol rate of 200 Ms/sec.

In order to provide more insight into the performance of the supported modulation and coding schemes in figure 7-3, simulated performance results are compared with channel capacities. The big solid blue curve is the unconstrained AWGN capacity from which the design of the desired set of spectral efficiencies has been started.

On the same plot, beside the performance results of the schemes (i.e., the required E_s/N_0 to achieve the target performance of $\text{BER}=10^{-6}$), which are reported with long dashed lines, two other theoretical limits are also reported. The solid curves are the *modulation-constrained capacities*, i.e., the theoretical limits over AWGN with the specific modulation schemes. Short dashed curves are the theoretical limits obtained using the Pragmatic approach, Bit Interleaved Code Modulation (BICM). This approach leads to a performance degradation that can be minimized using Gray labelling, but cannot be eliminated. The two theoretical limits are obtained in the limit of infinite codeword sizes. For a description of the methods used to derive these limits see references [B9] and [B10], and reference therein.

Comparing the performances of proposed schemes against the true limits (short dashed lines), one can observe that the loss of the coding schemes is uniform for all ACM formats and around one dB. Hence, the non-perfectly uniform spacing of curves in figure 7-1 is partly due to the choice of the modulation schemes, which have a variable loss against unconstrained capacity, and partly to the variable loss of the BICM approach. In particular, 8PSK and 32APSK have in general a bigger modulation loss, while the BICM approach, which gives almost negligible losses for 8PSK, 16APSK, and 64APSK gives a non-negligible loss for 32APSK modulation, which in fact does not admit Gray labelling.

Finally, although the performance curves are not uniformly spaced, the high density of the curves, and the partial overlapping between adjacent modulations, guarantees that the best trade-off can be found in all channel types and conditions, including the variable and non-linear satellite channel. In this latter case, in fact, the relative losses of the modulations are different so that curve of figure 7-1, relevant to the AWGN channel, will no longer be valid.

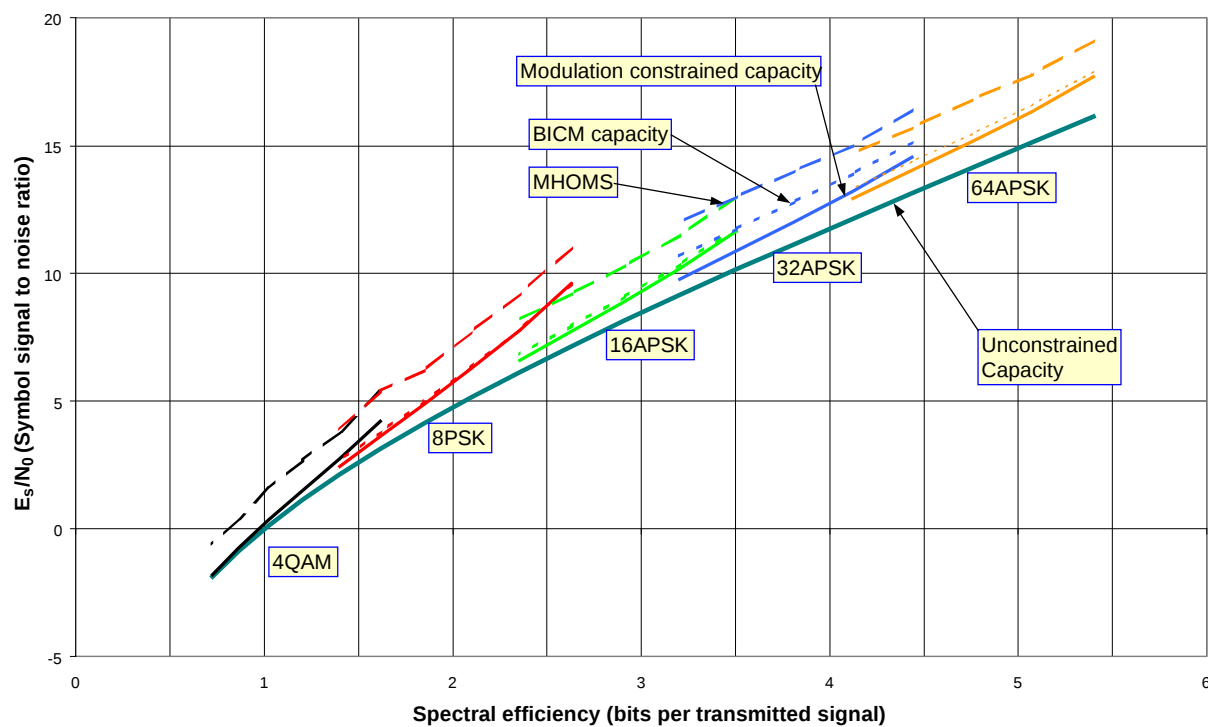


Figure 7-3: Comparison of the SCCC Achievable Performance versus Theoretical Bounds

ANNEX A

ACRONYMS

ACM	Adaptive Coding and Modulation
ASM	Attached SYNC Marker
AOS	Advanced Orbiting Systems
AWGN	Additive White Gaussian Noise
BER	Bit Error Ratio
CCSDS	Consultative Committee For Space Data Systems
FD	Frame Descriptor
FER	Frame Error Ratio
FM	Frame Marker
PL	Physical Layer
TC	Telecommand
TM	Telemetry
SCCC	Serially Concatenated Convolutional (Turbo) Code
VCM	Variable Coding and Modulation

ANNEX B

INFORMATIVE REFERENCES

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ANNEX C

INTERLEAVER DESCRIPTIONS

In order to support 27 distinct ACM formats, it is necessary to design only 19 permutations that allow a parallel implementation of the decoder with a degree of parallelism, $M=120$. Thus, the sizes of all 19 interleavers are integer multiples of 120. The outer encoder writes/reads its data in the memory in natural order. As for the inner encoder, the design of the permutations permits description of them in a simpler way than just giving the 19 series of integers from 0 to I .

The inner encoder writes/reads data according to the following relationship, which gives the writing/reading address at time i :

$$\pi(i) = W \times \left(i/W + \beta(i_w) \right)_M + \alpha(i_w) \quad i = 0, \dots, I \quad (\text{A-1})$$

where α and β are two vectors of length equal to $W=I/120$, and

$$i_w \equiv i \bmod W \quad (\text{A-2})$$

The elements of α (addresses of macrodata) range in $[0, W-1]$ and the elements of β (cyclic shifts of macrodata) range in $[0, 119]$.

Beyond yielding a more compact description of the interleaver, the previous permutation relationship allows for efficient implementation of a parallel decoder with parallelism 120 as shown in figure C-1.

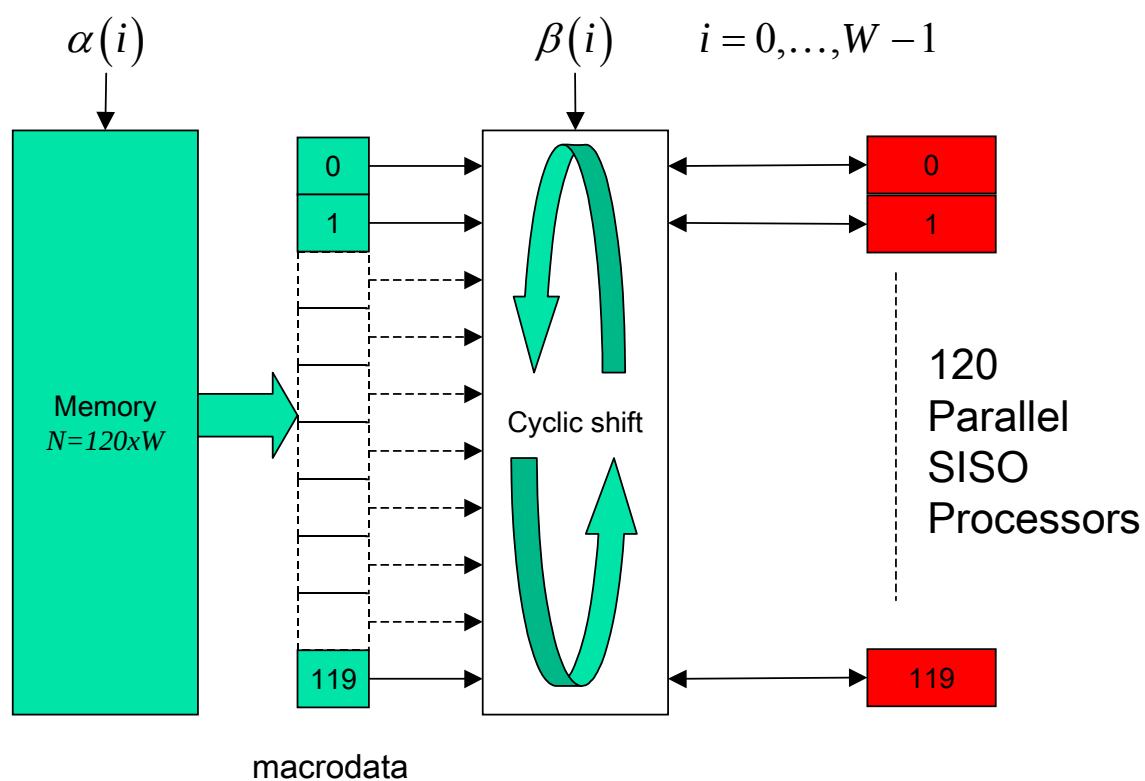


Figure C-1: Structure of the Simplified Interleaver

Table C-1: Interleaver Parameters (1-10)

	1		2		3		4		5		6		7		8		9		10	
	N=8640 W=72 S1=84 S2=84		N=10440 W=87 S1=90 S2=90		N=12600 W=105 S1=102 S2=102		N=14760 W=123 S1=110 S2=110		N=16920 W=141 S1=109 S2=109		N=19800 W=165 S1=127 S2=127		N=22320 W=186 S1=135 S2=135		N=25560 W=213 S1=144 S2=144		N=28800 W=240 S1=153 S2=153		N=32040 W=267 S1=161 S2=161	
	α	β	α	β	α	β	α	β	α	β	α	β	α	β	α	β	α	β	α	β
0	63	116	70	82	85	97	106	1	60	14	116	87	18	114	50	32	109	98	63	93
1	33	69	53	27	91	27	82	89	107	55	42	34	95	1	134	71	179	66	238	10
2	64	103	63	101	1	47	115	49	4	31	66	36	68	116	200	47	129	13	92	78
3	56	92	67	106	0	69	17	43	102	68	77	47	108	72	68	89	214	109	189	72
4	59	77	10	59	35	104	58	20	87	53	13	38	11	119	74	15	186	45	73	25
5	5	73	29	22	3	90	81	24	89	76	99	25	104	61	1	18	128	72	172	68
6	58	58	47	105	77	66	113	41	117	42	84	69	118	108	32	40	60	8	94	96
7	52	46	52	49	98	57	28	3	68	19	123	8	172	20	102	69	112	91	49	14
8	61	107	82	33	51	50	12	78	67	113	1	80	181	84	33	79	8	93	228	6
9	47	1	58	15	13	29	2	111	96	58	102	60	19	95	154	119	183	117	219	5
10	57	52	65	111	63	60	76	107	15	110	73	101	121	92	21	29	170	64	197	46
11	35	90	72	7	97	86	10	67	140	105	109	82	106	93	22	111	225	114	173	47
12	29	32	2	74	78	35	40	10	114	44	21	77	75	88	103	61	35	15	232	111
13	60	119	40	86	49	37	108	117	112	17	48	89	73	60	162	108	97	77	195	108
14	36	3	4	70	68	25	9	38	22	28	6	30	109	16	130	54	87	35	144	42
15	24	68	31	71	7	84	74	51	9	109	47	57	142	101	83	16	88	85	21	80
16	10	60	36	100	103	117	64	74	34	16	142	52	69	32	38	99	188	6	134	102
17	20	30	7	31	96	53	92	40	21	34	124	93	117	100	190	105	47	99	161	104
18	66	83	23	13	57	72	31	87	116	11	150	73	103	80	180	55	165	107	135	19
19	17	63	77	62	75	115	35	94	62	24	104	112	116	70	61	116	156	88	102	67
20	31	41	45	39	67	59	101	14	28	79	72	19	113	62	104	75	52	55	259	31
21	16	50	17	1	58	106	26	63	105	95	4	118	83	87	167	114	233	15	101	97
22	67	6	79	80	25	7	27	44	79	35	119	10	78	118	168	21	151	95	200	43
23	28	24	35	104	59	52	69	64	121	20	135	78	138	0	100	101	0	20	3	72
24	18	22	56	56	65	119	70	104	5	47	43	50	147	58	191	81	17	17	7	53
25	9	101	33	48	60	23	68	8	6	29	10	1	85	102	26	19	235	115	115	50
26	6	62	78	19	14	65	34	16	125	39	3	65	59	91	142	67	157	71	114	49
27	62	57	38	64	74	32	112	54	50	70	156	38	97	98	10	43	70	22	31	65
28	2	18	37	38	22	13	1	119	41	54	118	17	91	66	178	91	191	67	9	89
29	27	66	60	35	42	40	55	53	78	93	138	110	46	56	140	33	44	94	255	65
30	71	9	39	25	87	14	105	106	77	89	163	31	145	52	87	63	50	11	2	76
31	34	37	30	45	95	102	20	59	53	84	2	84	136	77	158	13	83	48	22	57
32	14	81	19	37	29	91	18	21	54	63	31	67	22	55	45	90	223	28	89	37
33	39	114	8	69	10	79	41	97	33	105	159	41	169	57	145	95	155	87	90	60
34	68	115	6	24	104	18	61	70	106	117	88	26	111	27	207	46	216	111	240	87
35	51	82	21	58	24	39	91	65	82	43	68	43	105	85	64	3	192	79	50	70
36	21	26	80	94	31	11	36	90	31	103	63	35	158	22	184	24	30	28	74	12
37	13	29	9	99	30	43	84	22	132	10	117	4	2	52	4	66	160	49	20	36
38	37	94	57	6	76	95	14	26	80	37	103	28	53	26	56	4	54	57	236	89
39	23	39	1	98	52	88	109	60	11	22	97	107	119	4	173	52	146	32	98	17
40	40	48	55	53	61	110	48	27	63	25	155	84	41	43	101	50	69	41	239	90
41	41	98	83	40	99	0	83	35	38	73	8	49	16	112	8	112	26	6	1	99
42	30	76	15	90	66	77	23	72	26	87	113	33	153	75	85	5	136	102	251	103
43	44	5	62	92	5	117	25	49	108	118	129	3	152	78	201	58	126	60	155	28
44	25	112	51	88	4	83	57	45	27	88	143	9	149	119	29	17	195	100	160	94
45	65	35	20	29	83	63	73	76	39	91	80	75	173	106	41	105	41	23	150	27
46	48	71	14	79	17	22	111	114	8	42	128	77	144	44	42	100	182	74	153	15
47	69	20	25	91	88	30	77	82	90	13	96	18	139	69	170	34	49	84	80	35
48	26	96	66	117	48	94	7	14	18	99	162	90	4	100	43	62	190	18	108	101
49	43	85	28	113	2	81	95	112	131	80	151	16	185	2	115	74	27	34	86	58
50	55	110	86	77	86	16	62	46	119	57	141	2	157	35	138	76	10	47	165	92
51	0	52	0	66	38	100	86	36	32	51	115	58	29	13	199	1	33	62	167	113
52	45	19	12	33	102	2	80	17	135	5	76	13	180	48	108	86	237	113	196	45
53	50	102	11	111	82	31	5	12	17	114	148	106	34	117	88	117	217	59	67	110
54	70	55	26	61	72	41	8	95	100	40	25	37	127	95	105	39	100	4	262	38
55	7	14	27	51	21	10	4	33	71	97	27	99	115	31	90	30	4	37	29	40
56	32	44	5	18	37	75	54	62	51	2	110	68	129	71	179	92	80	46	104	21
57	12	9	68	52	6	5	29	56	137	75	30	44	79	53	76	22	196	34	48	61
58	42	15	75	11	11	62	97	88	76	69	78	11	86	8	204	82	187	11	43	1

EXPERIMENTAL SPECIFICATION FOR FLEXIBLE SCC TURBO CODE FOR TELEMETRY APPLICATIONS

59	11	80	84	55	23	46	56	99	99	47	38	6	25	108	121	51	75	83	65	41
60	46	89	69	3	47	93	67	23	37	10	11	20	32	11	106	23	145	82	61	38
61	1	106	74	26	36	108	38	105	111	94	95	39	3	44	30	38	43	25	166	77
62	8	119	59	103	93	81	100	69	133	59	71	81	58	65	54	60	102	92	143	107
63	15	100	85	46	81	113	93	102	13	8	149	15	93	38	17	54	68	10	202	20
64	53	45	71	109	69	101	72	101	57	18	106	65	88	37	70	11	22	30	19	69
65	3	17	43	73	100	24	32	108	61	4	60	116	168	29	37	97	103	1	129	11
66	19	43	49	114	90	48	43	81	74	86	87	30	33	20	165	41	116	97	33	29
67	4	87	41	42	45	6	37	0	84	111	51	96	43	25	176	6	117	53	188	84
68	22	88	48	21	18	21	22	80	120	82	86	95	38	9	156	0	154	105	132	56
69	49	34	32	96	56	27	0	30	72	67	70	63	36	12	98	9	193	47	53	100
70	54	11	44	10	20	35	102	109	35	38	85	71	87	17	125	80	99	31	229	4
71	38	107	61	84	79	70	51	117	66	85	112	55	150	68	92	107	236	20	258	118
72			50	43	39	56	75	93	127	65	5	62	14	35	120	35	73	81	209	59
73			73	116	46	86	122	91	98	74	0	41	27	7	133	20	161	96	263	119
74			34	68	53	84	118	39	95	78	23	97	167	13	65	87	42	67	201	74
75			46	5	54	112	47	15	29	115	90	103	133	82	122	110	105	70	51	82
76			64	108	32	97	39	66	55	101	69	52	0	87	110	7	141	63	23	118
77			13	94	19	116	114	84	92	81	37	86	100	64	151	78	231	39	32	24
78			18	15	15	76	19	86	86	107	18	60	184	108	149	83	140	2	235	97
79			54	82	28	9	65	73	93	31	14	8	7	103	71	114	162	94	124	54
80			76	85	73	109	98	4	139	32	154	100	12	57	212	36	177	12	193	44
81			42	0	33	96	6	54	123	15	49	83	28	22	117	73	72	101	91	106
82			24	77	16	115	121	74	113	16	130	54	57	18	3	55	57	79	46	64
83			16	55	12	18	45	77	85	106	26	15	82	42	19	98	32	58	88	30
84			22	119	44	42	46	83	2	52	35	46	128	45	6	88	206	75	106	114
85			3	3	9	106	94	34	1	27	121	67	84	111	114	27	29	45	24	51
86			81	75	55	68	42	19	14	71	59	78	60	75	75	81	3	61	181	93
87					62	74	107	47	16	3	160	105	176	23	52	95	222	118	242	1
88					80	22	50	96	124	19	144	49	51	5	113	43	204	7	146	10
89					8	95	99	7	136	55	7	25	8	4	11	15	58	42	118	89
90					41	102	79	31	45	9	32	29	166	79	188	77	202	104	170	70
91					101	59	53	78	103	91	57	93	10	105	53	103	137	54	127	105
92					71	57	60	116	104	112	114	50	124	84	57	70	45	52	257	8
93					92	46	33	13	25	68	61	79	74	54	109	42	135	93	156	3
94					89	44	16	1	10	65	75	104	24	10	128	118	15	88	205	110
95					94	66	117	86	118	33	67	111	131	49	157	116	31	39	55	16
96					40	4	88	58	75	110	158	57	17	88	127	85	2	103	47	52
97					43	19	119	28	91	102	92	64	80	114	195	11	11	65	151	62
98					27	50	85	68	49	116	33	14	6	47	152	12	48	98	221	101
99					64	39	30	20	128	83	134	72	177	96	96	104	62	118	119	57
100					84	99	87	33	81	23	41	53	21	68	160	3	178	116	16	115
101					50	34	13	114	7	50	40	98	130	113	47	53	1	5	26	18
102					70	11	90	6	36	93	54	94	90	50	51	96	94	68	95	6
103					26	64	3	47	59	75	132	66	71	63	58	68	23	96	82	47
104					34	17	71	110	110	66	126	23	120	116	196	4	198	30	249	66
105							15	36	0	5	105	99	56	72	171	100	111	115	122	108
106							103	98	40	14	146	37	154	110	209	68	166	103	241	55
107							120	62	138	71	55	77	107	30	60	106	215	42	111	7
108							59	89	101	103	53	47	65	77	46	8	91	17	176	33
109							21	37	23	90	131	69	102	115	182	29	127	0	208	25
110							78	57	122	63	19	5	182	15	169	61	158	110	260	22
111							116	100	3	69	45	105	52	48	63	113	171	27	213	21
112							89	27	70	22	153	119	170	89	135	64	229	81	190	32
113							24	113	48	108	93	73	171	25	194	70	159	23	131	13
114							49	10	52	119	34	82	89	39	192	32	221	1	72	48
115							66	32	109	35	20	19	55	46	118	58	67	49	149	80
116							44	103	88	88	12	36	114	97	69	77	85	109	18	84
117							63	5	44	28	74	114	81	21	111	10	194	73	35	4
118							11	97	69	117	111	41	26	28	7	116	77	72	252	82
119							104	81	24	0	157	75	92	14	147	59	65	38	10	55
120							52	59	46	82	36	110	183	105	153	40	150	21	139	39
121							96	56	126	1	62	45	1	29	136	45	153	119	141	53
122							110	99	19	21	17	109	49	74	2	3	142	33	110	65
123									115	113	120	86	35	110	187	30	9	59	158	75
124									47	53	98	117	63	34	25	52	149	45	100	103

EXPERIMENTAL SPECIFICATION FOR FLEXIBLE SCC TURBO CODE FOR TELEMETRY APPLICATIONS

125										20	86	140	0	163	40	210	90	181	61	168	87
126										42	7	91	84	9	99	163	28	180	56	223	100
127										65	73	15	28	179	90	23	115	19	66	217	26
128										83	46	58	112	143	66	80	49	173	40	198	71
129										64	95	122	31	50	93	12	16	56	7	222	64
130										56	47	56	80	137	19	112	1	20	73	99	81
131										134	93	83	91	125	20	91	119	95	47	17	112
132										130	39	94	6	135	118	131	82	40	89	133	79
133										73	100	127	21	20	82	84	111	234	86	145	96
134										129	86	152	22	62	76	155	60	197	50	59	104
135										58	105	145	12	40	23	185	62	218	3	112	5
136										30	79	161	25	31	83	132	44	81	69	116	90
137										94	2	89	118	72	107	97	79	219	85	44	45
138										43	61	137	98	174	31	20	39	139	19	246	42
139										12	98	65	24	148	10	144	71	205	91	37	25
140										97	99	64	56	175	77	16	56	106	113	70	19
141												24	65	98	55	77	109	213	13	187	12
142												108	8	162	65	86	26	118	24	45	113
143												164	68	47	95	175	9	133	6	120	63
144												82	89	156	91	119	99	209	32	140	29
145												79	42	64	1	0	9	119	90	206	36
146												100	97	110	12	148	105	189	52	138	14
147												101	2	30	52	55	0	55	29	66	59
148												29	103	101	26	164	107	13	110	137	8
149												22	59	132	43	39	34	59	107	211	99
150												133	82	48	15	139	108	207	22	233	91
151												107	46	178	37	49	66	64	37	41	34
152												136	115	5	17	67	57	122	84	256	104
153												52	107	77	36	31	47	152	15	162	51
154												81	106	141	7	211	74	28	93	216	30
155												147	95	155	42	161	23	132	64	38	86
156												39	55	122	6	28	25	174	111	56	74
157												139	33	37	112	78	32	210	9	34	107
158												50	17	23	51	197	50	92	44	36	85
159												44	71	13	116	14	102	232	0	185	116
160												16	72	42	49	198	104	208	14	182	9
161												125	60	44	98	141	112	224	70	234	78
162												9	64	70	79	206	42	96	95	212	2
163												28	74	126	108	40	28	61	80	93	119
164												46	101	96	106	66	83	79	114	175	67
165														151	54	193	69	125	25	54	42
166														39	40	15	27	114	48	107	76
167														140	8	35	94	101	112	226	31
168														76	90	183	95	76	62	220	95
169														165	38	82	97	86	3	87	44
170														61	35	166	91	203	78	64	17
171														45	64	73	29	108	16	97	109
172														123	13	205	5	238	53	142	55
173														54	31	202	21	84	63	79	49
174														164	47	48	50	147	8	245	37
175														94	80	189	113	134	57	25	35
176														67	24	79	15	25	102	121	115
177														134	74	174	37	63	22	125	83
178														99	11	24	35	14	75	148	58
179														146	57	34	59	37	0	126	118
180														112	60	36	42	201	38	244	98
181														66	69	181	63	124	106	207	0
182														160	9	159	103	239	89	230	68
183														15	85	59	14	212	105	40	78
184														159	18	13	40	74	13	0	28
185														161	75	186	73	16	20	8	57
186																129	89	228	18	250	20
187																123	38	36	91	210	35
188																94	93	167	66	11	61
189																124	65	143	34	62	22
190																203	86	89	26	83	77

EXPERIMENTAL SPECIFICATION FOR FLEXIBLE SCC TURBO CODE FOR TELEMETRY APPLICATIONS

191																99	2	5	50	224	18
192																9	21	121	11	194	52
193																126	67	184	69	237	93
194																89	13	138	100	27	116
195																95	54	211	97	58	80
196																93	116	164	77	5	47
197																172	31	144	58	171	106
198																208	85	110	79	179	117
199																107	46	104	104	261	49
200																27	85	226	98	57	0
201																18	81	24	6	77	73
202																72	7	131	43	6	82
203																5	11	38	76	113	112
204																150	98	168	71	105	84
205																81	84	34	60	183	89
206																116	3	230	45	85	101
207																177	119	51	74	96	95
208																143	114	71	28	227	45
209																146	110	18	32	84	27
210																62	58	113	117	253	15
211																137	27	123	59	204	10
212																44	45	200	44	218	113
213																		21	55	109	110
214																		227	83	254	59
215																		53	81	117	4
216																		148	4	248	28
217																		172	51	152	72
218																		199	113	71	91
219																		107	40	69	3
220																		115	41	68	32
221																		82	97	78	75
222																		185	10	215	94
223																		93	1	12	103
224																		130	17	184	17
225																		120	31	178	24
226																		46	64	225	40
227																		12	27	180	41
228																		163	30	266	71
229																		176	102	186	61
230																		169	107	264	13
231																		98	39	199	65
232																		7	25	265	86
233																		39	77	164	107
234																		66	105	174	54
235																		175	87	76	23
236																		90	52	13	5
237																		6	57	14	102
238																		78	15	28	9
239																		220	64	243	108
240																				191	19
241																				136	60
242																				231	33
243																				30	63
244																				130	92
245																				128	30
246																				60	56
247																				39	106
248																				103	45
249																				214	63
250																				81	6
251																				123	64
252																				15	117
253																				154	1
254																				75	11
255																				192	74
256																				159	25

EXPERIMENTAL SPECIFICATION FOR FLEXIBLE SCC TURBO CODE FOR TELEMETRY APPLICATIONS

257																				163	119
258																				247	3
259																				52	88
260																				147	37
261																				203	42
262																				169	70
263																				157	82
264																				177	100
265																				42	118
266																				4	31

Table C-2: Interleaver Parameters (11 to 19)

	11		12		13		14		15		16		17		18		19	
	N=35280 W=294 S1=158 S2=158		N=38880 W=324 S1=166 S2=166		N=42480 W=354 S1=174 S2=174		N=46440 W=387 S1=193 S2=193		N=50040 W=417 S1=189 S2=189		N=54000 W=450 S1=208 S2=208		N=57960 W=483 S1=215 S2=215		N=61560 W=513 S1=222 S2=222		N=65520 W=546 S1=216 S2=216	
	α	β	α	β	α	β	α	β	α	β	α	β	α	β	α	β	α	β
0	30	28	193	72	238	48	374	11	147	48	211	24	153	110	380	92	282	83
1	165	81	208	10	260	88	269	31	76	100	353	68	307	101	415	60	437	63
2	156	103	3	107	253	11	21	35	22	84	430	37	274	42	81	60	397	89
3	235	46	254	13	208	43	84	46	64	92	56	55	291	45	327	41	471	115
4	15	75	165	59	130	102	6	45	53	64	242	50	158	50	60	21	218	29
5	126	82	81	50	190	79	91	85	50	39	440	76	141	75	289	28	62	79
6	203	34	154	42	193	2	81	33	219	63	111	67	241	112	341	16	65	0
7	270	112	38	77	122	73	227	115	134	16	354	11	80	62	162	31	396	88
8	189	71	104	83	118	47	324	49	406	0	123	96	310	1	409	104	172	93
9	103	85	216	52	317	95	367	36	315	118	109	89	83	66	95	116	182	49
10	145	1	63	68	177	42	225	64	416	62	80	23	110	82	159	29	84	88
11	186	92	68	109	224	111	294	108	187	79	126	63	433	105	330	105	116	48
12	264	7	116	4	143	89	159	28	307	1	396	51	17	47	262	57	120	98
13	218	109	292	44	181	100	277	110	252	85	51	106	386	18	157	9	224	39
14	52	57	31	13	178	86	265	27	373	106	34	2	298	38	223	115	0	41
15	265	55	167	57	29	30	87	34	146	104	302	73	89	59	170	63	174	107
16	141	24	322	114	64	40	193	53	115	87	386	25	478	54	34	118	424	0
17	181	106	94	65	173	81	258	15	293	29	161	94	40	105	137	44	362	105
18	208	108	122	67	315	50	8	56	291	10	292	79	279	63	411	25	247	52
19	9	51	21	19	198	3	141	48	33	97	30	16	361	60	269	13	511	82
20	88	76	201	90	46	7	210	61	257	75	267	8	130	41	174	110	343	95
21	196	100	86	0	258	0	89	4	399	90	163	21	230	99	197	103	245	59
22	32	19	23	24	108	98	205	23	213	94	337	6	425	12	4	64	15	105
23	78	40	270	15	106	69	195	67	4	46	202	5	35	22	202	37	262	96
24	130	17	236	78	7	50	279	33	207	43	147	54	19	95	110	49	300	60
25	162	29	7	36	167	101	338	118	233	70	232	49	280	102	437	58	207	11
26	269	83	227	43	308	51	82	98	266	6	420	2	166	37	180	7	26	31
27	42	118	307	87	139	105	151	59	39	25	426	14	49	76	375	94	130	35
28	99	26	316	63	68	85	202	10	249	116	221	110	248	25	368	3	539	46
29	22	74	44	78	262	96	40	40	405	51	28	83	119	17	337	55	230	45
30	192	31	27	21	125	19	63	69	226	66	437	115	71	38	475	42	136	85
31	147	80	297	32	246	91	90	78	341	17	262	77	170	13	303	81	305	33
32	221	3	285	85	129	114	358	94	203	98	324	40	466	76	176	23	281	36
33	16	38	279	21	295	30	10	80	73	73	20	43	387	95	320	79	285	64
34	182	77	175	119	87	119	121	94	110	34	307	69	123	6	397	36	485	108
35	188	107	192	70	72	17	85	5	62	26	1	85	327	36	169	62	209	28
36	73	11	37	103	332	9	158	79	89	10	37	111	200	51	447	73	145	110
37	159	10	293	92	226	1	318	76	359	15	77	105	4	34	372	77	500	27
38	256	48	239	80	232	113	16	73	204	81	359	38	333	16	291	84	67	34
39	21	2	32	96	94	13	167	50	58	15	3	32	464	13	134	38	2	53
40	10	63	162	118	307	22	322	88	124	93	264	112	103	46	84	66	436	15
41	72	102	30	114	67	109	68	1	358	18	412	94	142	39	443	96	491	56
42	61	59	310	38	329	24	168	24	247	102	410	52	470	112	154	34	512	48
43	169	52	218	30	98	6	119	114	104	50	435	108	108	71	385	26	277	61
44	232	41	209	49	282	94	303	1	250	88	185	88	67	70	190	111	232	4
45	283	54	146	94	303	70	263	37	116	76	261	7	93	24	51	81	259	23
46	271	85	315	68	142	67	28	31	323	119	79	99	253	111	423	102	355	67
47	6	66	2	82	55	32	350	20	152	65	220	27	315	52	500	47	310	118
48	171	86	206	12	1	79	172	2	170	19	206	70	451	83	165	19	384	98
49	29	12	265	62	250	90	382	19	362	50	175	117	105	56	501	0	261	10
50	279	12	281	110	10	86	174	18	156	71	168	20	269	84	357	29	265	40
51	115	61	191	31	318	80	293	47	237	11	186	48	442	42	469	71	367	69
52	199	8	184	54	274	23	114	118	121	30	263	111	175	97	226	70	513	94
53	285	93	282	6	306	99	30	27	238	27	394	66	155	68	377	65	475	80
54	231	67	305	99	233	52	234	91	285	23	280	72	51	73	151	22	228	94
55	272	114	9	97	284	89	333	112	385	76	38	31	33	74	482	53	184	5
56	194	5	140	34	21	29	96	102	240	73	49	104	216	80	358	5	323	76
57	129	27	214	48	255	110	247	51	65	9	289	96	271	64	206	39	44	73
58	210	98	158	60	289	69	215	14	95	21	407	45	2	45	12	86	398	50

EXPERIMENTAL SPECIFICATION FOR FLEXIBLE SCC TURBO CODE FOR TELEMETRY APPLICATIONS

59	19	89	109	85	163	18	189	72	175	53	365	19	357	90	294	15	258	111
60	253	115	278	103	6	58	380	8	316	41	245	107	448	77	192	71	221	1
61	163	64	289	8	343	107	252	45	348	99	447	119	101	9	381	9	446	24
62	263	69	42	12	144	118	297	70	374	78	23	37	403	82	346	113	141	114
63	255	113	126	40	124	27	307	92	44	6	223	51	432	79	8	2	444	1
64	259	99	55	95	96	72	329	98	93	119	103	42	370	43	384	40	201	37
65	275	45	74	58	221	87	117	101	60	116	290	46	356	47	243	12	216	31
66	282	87	274	22	188	46	102	25	300	47	443	93	203	23	379	23	39	20
67	175	96	242	29	111	38	154	107	334	22	296	99	339	22	394	112	45	15
68	83	117	13	42	323	20	250	5	310	44	224	9	249	21	344	98	492	2
69	227	47	58	117	172	107	136	54	88	60	266	10	435	10	237	4	203	19
70	288	19	129	26	214	65	238	17	162	108	73	114	106	3	486	86	168	18
71	164	53	143	84	169	84	330	32	212	28	201	17	380	7	251	72	227	47
72	108	42	323	4	89	15	199	3	41	72	193	44	383	0	147	117	1	23
73	202	104	114	110	247	93	50	70	21	82	130	97	56	107	193	27	284	113
74	82	25	34	41	297	62	363	52	28	49	284	22	346	29	326	80	64	27
75	55	101	69	27	342	4	266	9	264	58	120	41	382	34	119	107	135	91
76	3	93	204	36	270	64	161	77	57	47	195	2	300	30	462	51	211	112
77	26	36	59	15	286	103	45	63	317	49	214	36	469	4	488	111	435	4
78	89	41	43	66	335	102	245	34	254	4	327	18	254	14	200	100	17	10
79	140	51	229	20	9	37	133	74	392	56	100	118	364	89	400	24	443	102
80	20	39	77	37	325	115	126	44	10	55	136	56	94	86	124	91	179	51
81	178	30	33	10	346	111	83	19	137	20	87	86	456	31	1	101	322	14
82	40	37	179	86	191	45	116	49	239	87	196	87	95	104	455	6	369	72
83	291	29	148	45	293	10	185	22	286	7	275	60	330	94	481	62	522	8
84	53	23	159	79	34	44	286	40	2	81	346	24	338	20	338	97	479	45
85	293	15	117	106	121	94	346	64	26	12	338	26	304	110	145	82	393	70
86	17	106	212	98	99	0	152	7	382	114	349	90	471	32	207	53	417	92
87	28	6	152	97	264	14	53	62	377	108	167	82	181	114	29	85	473	101
88	209	58	98	25	334	47	64	83	283	39	425	91	266	115	86	56	504	25
89	251	84	151	24	86	113	7	92	190	1	155	101	440	96	205	54	478	107
90	176	22	160	64	185	7	244	68	282	83	105	115	388	71	94	93	53	54
91	289	68	182	100	136	80	135	55	172	67	119	69	452	1	487	8	123	17
92	134	7	103	61	338	104	257	116	353	32	41	61	135	8	21	57	332	32
93	65	0	48	98	14	92	148	60	183	2	241	113	384	3	26	5	453	119
94	243	28	120	23	39	33	26	17	159	14	95	71	428	5	224	99	195	9
95	206	63	238	17	339	40	97	26	9	98	306	64	404	55	39	84	337	77
96	154	19	28	91	107	25	22	115	389	116	422	59	447	98	97	14	92	63
97	56	77	180	87	61	59	315	93	206	17	371	0	190	35	58	24	423	73
98	27	110	121	72	160	48	345	28	67	112	279	34	154	42	234	60	271	34
99	184	111	221	75	90	70	253	85	265	89	125	95	294	2	214	76	374	3
100	94	73	101	33	187	32	232	111	274	5	234	76	245	109	222	87	421	74
101	75	3	64	71	78	31	184	104	378	86	117	80	399	117	392	110	304	44
102	119	13	147	63	230	33	77	66	325	110	364	62	368	92	416	114	335	93
103	69	91	135	55	73	36	143	88	98	68	21	30	134	26	72	26	5	19
104	35	81	72	74	352	2	251	11	78	43	377	78	431	108	188	18	406	49
105	284	118	115	99	104	78	98	6	227	109	177	57	244	48	14	31	353	115
106	127	90	138	82	305	16	260	96	102	114	332	13	11	27	257	51	75	22
107	250	97	269	1	340	105	4	76	281	52	392	74	136	31	454	39	484	23
108	123	46	12	73	197	53	275	58	410	12	169	53	29	111	352	19	52	64
109	260	9	82	52	265	17	278	39	109	78	331	16	179	54	2	39	292	7
110	143	44	52	108	84	64	49	75	267	115	151	58	343	65	178	58	103	62
111	240	116	210	101	174	63	353	53	368	91	311	75	208	95	85	98	100	92
112	204	74	19	50	299	8	372	71	355	36	29	51	233	116	458	115	541	68
113	85	114	166	35	80	95	366	106	143	91	384	61	221	49	493	7	450	55
114	225	0	189	18	312	15	180	73	14	62	129	14	446	111	351	31	291	116
115	274	101	57	92	37	102	240	90	361	38	446	27	318	44	32	102	153	89
116	70	82	0	87	202	119	370	3	222	34	181	15	402	100	184	108	386	26
117	36	22	222	112	320	75	343	113	196	106	165	90	250	33	264	75	339	35
118	0	35	290	53	278	25	273	80	356	105	428	21	441	80	343	107	41	118
119	38	31	257	83	152	116	37	52	241	48	256	105	31	79	23	45	519	64
120	112	50	194	116	47	34	298	97	129	99	191	29	415	39	417	59	48	4
121	290	42	157	22	182	22	368	57	184	35	237	109	129	103	201	65	505	28
122	179	76	237	46	231	44	235	12	108	74	114	18	389	99	229	68	21	111
123	233	94	235	28	313	59	111	43	42	95	131	32	211	76	474	119	390	104
124	128	21	111	53	341	86	128	103	197	55	27	40	240	87	24	89	464	66

EXPERIMENTAL SPECIFICATION FOR FLEXIBLE SCC TURBO CODE FOR TELEMETRY APPLICATIONS

125	8	27	168	91	162	106	147	37	100	105	313	4	287	67	16	51	415	10
126	247	70	299	5	301	67	213	109	105	61	116	52	68	108	426	89	410	11
127	139	15	61	30	123	115	268	100	302	117	198	68	82	93	102	92	447	6
128	193	39	173	68	281	13	288	4	23	101	68	78	270	75	50	3	480	96
129	44	34	131	16	259	97	311	29	371	103	406	8	276	104	203	67	29	76
130	220	55	1	14	175	20	103	56	133	46	411	63	24	18	150	77	77	58
131	216	103	155	41	133	75	108	81	191	107	413	53	138	0	80	25	76	39
132	60	116	56	115	183	57	59	28	31	83	423	69	302	62	464	18	375	19
133	183	85	88	77	33	91	57	32	388	26	137	45	457	93	370	64	442	47
134	80	32	108	62	235	66	226	105	364	13	121	81	354	8	44	73	106	68
135	180	40	6	38	3	99	149	33	260	111	351	54	413	74	212	8	246	24
136	239	77	217	96	135	104	187	38	142	102	281	86	234	88	45	11	237	21
137	168	62	301	34	113	26	46	71	397	31	238	66	482	48	166	33	133	77
138	114	54	145	88	119	50	386	7	179	31	45	112	422	70	195	6	131	16
139	14	79	150	10	348	1	352	54	218	77	350	95	267	105	82	47	150	14
140	120	18	317	30	314	11	373	117	186	45	300	71	220	28	319	50	349	53
141	205	65	286	117	213	35	204	16	154	118	180	28	20	116	442	82	427	59
142	131	87	90	81	292	42	337	35	12	80	378	82	437	53	108	61	461	106
143	117	105	207	76	88	4	203	78	120	69	13	119	151	101	354	2	445	62
144	190	33	304	77	12	52	101	0	25	48	219	91	353	114	191	81	331	90
145	93	71	65	107	83	21	179	30	393	85	269	94	14	98	349	34	118	3
146	261	13	321	118	52	83	312	67	208	93	53	33	10	17	322	78	33	113
147	62	83	188	69	345	6	134	23	313	57	312	106	303	107	52	95	303	80
148	174	2	267	60	257	95	376	81	298	0	259	102	187	63	348	14	527	97
149	170	14	36	67	242	78	281	61	201	24	277	114	150	117	181	21	159	57
150	172	59	137	11	271	85	296	59	125	66	176	116	157	4	18	62	229	12
151	34	45	54	104	69	62	310	44	366	87	204	98	472	24	406	48	42	90
152	2	7	246	93	205	19	201	62	119	4	385	104	137	85	158	52	320	43
153	276	24	17	101	290	31	61	57	18	54	375	73	149	113	156	36	455	86
154	64	88	123	39	249	114	264	99	193	97	200	1	113	65	270	104	112	103
155	74	60	224	40	298	88	160	86	396	88	326	81	460	18	471	103	414	65
156	107	115	240	47	40	66	130	119	66	84	212	93	410	118	399	54	326	37
157	213	38	260	119	140	77	236	21	27	100	194	3	121	89	350	106	405	109
158	102	70	252	7	272	109	347	14	221	104	32	12	90	69	49	106	68	100
159	101	66	215	95	330	46	70	91	171	18	4	24	237	18	142	35	382	29
160	155	113	241	50	210	96	218	41	13	35	438	83	391	57	277	17	87	59
161	201	102	8	32	154	56	359	95	48	67	439	49	263	77	292	92	104	94
162	76	49	313	18	328	53	267	26	277	21	108	72	193	43	196	105	97	70
163	158	106	113	51	217	82	207	46	72	41	7	57	374	35	388	46	219	56
164	110	29	288	36	227	118	357	48	200	23	88	108	50	50	89	15	180	6
165	13	8	213	73	201	40	76	80	319	30	62	19	88	84	408	95	196	81
166	226	36	102	111	291	79	65	99	311	74	72	49	183	92	434	15	399	94
167	195	89	169	102	141	98	290	77	99	8	276	55	225	58	268	30	152	105
168	258	108	190	78	236	101	93	95	161	26	115	25	224	3	28	114	528	16
169	161	101	144	9	74	5	384	15	395	43	282	88	202	40	90	79	146	33
170	48	104	133	66	126	29	29	93	299	25	16	21	449	14	230	109	361	71
171	79	111	243	94	77	61	239	50	131	40	187	77	81	64	220	20	253	22
172	71	92	248	23	27	13	14	12	243	16	369	17	73	7	140	48	125	40
173	25	5	219	104	13	8	78	42	8	103	36	6	63	106	46	94	86	36
174	121	75	71	86	149	43	274	114	74	90	99	27	406	54	227	28	489	54
175	217	95	127	44	54	90	115	9	391	107	363	85	275	37	453	43	336	117
176	24	112	256	2	245	99	328	73	357	96	133	73	296	69	87	1	38	44
177	167	35	29	60	43	81	80	79	68	51	142	13	344	11	164	118	99	28
178	135	97	250	63	326	37	15	88	34	20	255	43	436	46	121	86	391	38
179	207	23	47	100	19	42	35	84	91	38	47	26	308	50	483	5	34	75
180	54	26	263	55	304	76	378	17	16	0	178	96	18	87	396	27	240	78
181	238	83	320	87	311	110	336	33	37	28	251	67	377	41	283	90	206	0
182	257	64	141	19	8	45	344	111	318	92	403	80	143	16	275	117	70	30
183	46	19	202	3	331	23	254	13	141	3	67	11	163	82	187	16	387	39
184	223	96	262	57	36	65	211	32	69	52	14	7	111	11	309	32	154	61
185	173	48	309	51	120	71	340	102	232	113	408	105	421	103	37	69	161	73
186	266	107	273	79	116	55	20	77	407	33	149	22	184	57	221	102	169	97
187	228	17	302	65	155	34	223	1	415	1	301	12	169	119	61	59	191	26
188	31	38	70	46	283	21	349	82	97	64	154	39	76	13	422	116	23	91
189	236	20	67	59	11	69	5	106	169	22	424	68	177	12	248	24	299	99
190	241	1	76	93	333	29	129	70	176	13	172	60	43	1	361	96	28	119

EXPERIMENTAL SPECIFICATION FOR FLEXIBLE SCC TURBO CODE FOR TELEMETRY APPLICATIONS

191	98	37	174	108	23	94	42	3	127	49	401	34	265	55	284	11	408	100
192	67	4	66	94	240	84	150	27	402	64	343	14	132	22	314	62	403	91
193	23	56	130	117	156	88	19	97	220	95	50	88	78	91	6	42	537	41
194	224	68	271	113	254	9	71	2	61	111	26	89	236	81	256	69	223	49
195	90	119	259	84	49	116	316	108	261	61	145	74	331	59	62	23	185	95
196	292	50	20	112	252	94	332	38	280	90	303	101	272	79	366	38	205	27
197	214	26	15	40	50	46	56	41	87	75	18	47	195	74	38	18	117	38
198	142	41	183	106	209	83	191	75	198	110	48	82	334	51	238	93	27	78
199	104	30	253	70	319	74	27	67	242	71	8	103	222	34	261	41	544	115
200	211	10	87	12	207	28	220	94	246	19	308	15	176	111	389	63	198	52
201	246	16	223	32	277	0	34	113	303	17	429	79	116	44	508	67	178	46
202	96	100	296	61	158	93	73	114	130	98	5	101	194	20	198	97	202	48
203	267	71	284	68	62	20	173	35	335	70	310	51	32	53	17	53	22	72
204	278	22	153	0	239	117	32	13	333	14	383	103	259	97	78	54	139	45
205	41	10	303	14	322	71	54	59	258	63	183	64	371	91	359	58	431	95
206	12	48	89	18	164	59	381	49	70	70	270	58	159	102	40	58	338	15
207	219	67	99	56	302	57	23	61	312	106	101	44	400	27	209	85	186	50
208	252	60	40	26	171	12	249	63	49	45	333	36	396	45	304	42	503	12
209	148	57	318	98	138	100	319	31	210	33	321	37	405	109	213	40	346	42
210	150	98	187	29	35	108	51	98	188	86	320	33	206	118	100	100	458	114
211	273	5	125	31	18	23	361	22	367	54	421	44	423	116	436	66	378	9
212	50	47	25	68	75	103	283	25	287	67	416	97	465	85	274	114	394	79
213	87	107	308	20	176	68	79	116	324	73	271	56	337	9	280	44	4	21
214	152	58	276	24	324	27	206	83	177	29	118	59	393	119	175	43	496	88
215	92	96	26	5	241	69	379	107	400	65	389	117	463	38	98	75	439	82
216	280	90	170	1	206	47	228	101	338	78	82	94	345	31	329	10	272	107
217	151	84	171	37	56	57	112	117	19	10	203	84	351	73	296	45	364	5
218	215	118	272	82	321	49	292	23	36	37	295	89	395	68	115	2	156	66
219	191	27	93	85	223	1	156	96	151	101	444	7	332	84	66	110	315	1
220	37	118	78	2	263	63	383	42	81	63	83	34	147	41	129	115	521	4
221	249	15	247	88	92	111	224	47	174	88	344	57	189	36	92	55	389	84
222	106	93	91	4	273	38	256	81	163	112	75	67	109	30	218	116	236	76
223	116	79	60	72	109	18	295	79	383	112	297	65	8	38	339	99	538	33
224	245	62	203	8	300	77	242	72	269	50	405	31	127	75	143	84	66	104
225	146	8	228	5	30	97	280	3	387	10	230	41	104	83	69	65	3	32
226	149	114	181	74	91	33	216	24	215	83	97	20	107	62	74	70	233	102
227	51	87	39	63	76	112	74	107	106	27	102	4	0	95	278	22	193	41
228	109	0	255	43	296	81	186	89	123	57	210	10	91	109	318	12	278	30
229	286	72	142	98	216	51	24	110	290	76	106	70	325	66	313	29	407	13
230	58	15	14	84	251	104	88	40	209	99	218	16	148	98	460	107	190	82
231	125	34	97	76	146	45	110	48	262	51	361	29	278	42	266	56	95	79
232	100	22	4	47	347	72	122	109	332	82	379	92	84	37	267	26	329	106
233	45	95	176	90	4	114	259	0	0	43	273	115	167	24	127	111	263	70
234	68	110	110	54	180	60	335	51	394	29	441	109	292	71	428	55	462	27
235	18	115	124	83	93	37	331	110	107	96	85	107	229	106	105	37	366	17
236	1	12	112	13	66	22	255	56	347	56	35	109	398	78	103	72	404	48
237	230	80	51	16	16	35	123	53	29	31	46	63	6	105	63	28	94	2
238	248	73	95	36	170	73	86	54	273	116	305	116	408	28	68	88	357	108
239	43	117	132	101	179	23	198	113	211	44	40	1	30	101	449	33	122	55
240	33	94	107	14	285	48	229	6	45	106	135	30	72	56	387	21	330	119
241	198	19	275	105	24	85	67	64	253	15	9	76	61	21	323	76	307	67
242	200	92	258	66	70	40	169	34	86	23	434	45	227	112	31	46	392	14
243	157	103	83	118	112	84	105	87	158	56	427	16	54	32	300	101	456	57
244	47	54	156	25	248	4	306	10	5	5	397	52	192	15	504	11	314	31
245	86	1	196	77	192	26	75	31	292	62	93	54	246	60	302	77	252	32
246	95	64	85	58	168	58	157	8	413	71	228	42	160	72	503	19	448	35
247	4	46	149	75	15	59	200	91	330	75	184	69	480	3	93	78	289	86
248	237	59	312	44	97	117	127	17	234	9	322	23	27	31	41	63	302	83
249	97	55	198	33	349	43	214	76	320	37	414	26	381	102	245	9	486	37
250	81	61	294	102	229	108	371	69	136	32	146	92	36	66	363	25	242	59
251	262	52	41	64	128	107	313	88	122	36	366	72	64	115	287	57	222	39
252	113	68	75	35	211	90	118	22	52	77	250	87	434	110	232	13	400	69
253	118	31	277	21	350	15	351	9	351	39	268	19	164	61	473	102	162	63
254	84	89	84	73	145	8	137	85	217	103	209	26	455	69	83	32	50	109
255	63	36	199	99	196	14	3	5	379	81	314	112	261	23	478	34	457	51
256	138	24	53	41	150	109	47	104	343	72	69	48	38	14	459	73	517	98

EXPERIMENTAL SPECIFICATION FOR FLEXIBLE SCC TURBO CODE FOR TELEMETRY APPLICATIONS

257	136	39	73	3	151	42	36	14	327	41	395	76	329	64	374	94	269	104
258	7	66	139	80	65	63	334	98	354	35	246	0	336	19	35	93	248	93
259	137	51	298	110	131	96	171	12	308	46	319	3	439	43	163	38	24	25
260	287	112	136	70	115	3	181	102	248	64	182	62	324	21	391	97	143	112
261	122	5	226	89	25	48	131	20	173	109	10	96	348	7	429	70	465	2
262	229	40	232	48	215	6	72	101	145	7	225	95	481	22	376	39	110	23
263	242	41	118	34	161	67	217	55	344	89	153	50	28	51	148	113	301	103
264	197	44	319	27	195	10	289	78	255	79	402	77	185	90	308	87	148	101
265	49	33	177	109	101	31	272	36	372	117	104	117	427	89	425	23	509	117
266	277	87	62	91	48	17	270	60	43	58	94	102	409	107	362	83	251	96
267	66	59	233	64	17	27	285	106	79	91	404	99	443	25	450	108	30	84
268	144	46	266	114	51	54	177	52	7	21	134	24	426	67	210	31	20	1
269	160	71	195	15	316	78	39	66	349	118	373	114	289	70	301	3	370	105
270	59	21	230	85	186	65	48	65	231	38	387	90	444	60	107	17	372	64
271	57	17	280	95	219	64	106	21	205	74	64	23	125	76	445	81	8	16
272	212	76	185	20	71	1	276	118	59	22	139	104	42	96	265	7	510	21
273	105	38	100	17	280	106	163	57	360	5	358	38	92	10	342	71	363	81
274	234	0	24	71	165	36	369	34	245	6	442	54	350	52	120	76	126	31
275	5	30	128	22	184	76	317	62	309	24	316	25	379	72	254	49	296	40
276	166	91	163	94	189	15	302	7	294	45	150	86	173	107	410	60	514	107
277	177	50	225	72	327	28	69	92	301	58	388	24	218	39	355	28	347	0
278	254	97	105	114	132	62	43	105	150	34	12	42	311	82	440	98	12	73
279	153	70	96	87	294	32	183	45	138	54	362	96	256	93	36	118	469	110
280	185	88	119	57	85	77	348	74	268	105	158	79	235	45	364	16	467	26
281	132	111	92	6	105	2	2	19	275	8	156	6	133	80	413	69	55	51
282	133	10	134	115	269	20	209	112	297	102	86	38	342	26	7	21	516	63
283	268	23	22	102	38	68	208	108	148	114	61	43	140	49	324	18	93	9
284	11	106	10	113	127	118	212	43	295	1	70	15	420	6	217	79	220	72
285	91	63	244	59	22	95	327	116	272	93	173	8	397	90	161	27	19	56
286	187	69	200	46	275	50	230	107	363	113	57	87	168	48	505	12	102	12
287	111	7	172	93	288	54	120	69	412	18	144	78	335	12	402	61	268	61
288	244	14	306	28	204	95	139	97	322	22	170	61	16	78	490	4	238	47
289	124	90	231	41	81	53	341	40	103	0	376	2	392	80	340	105	508	53
290	281	21	283	92	28	64	304	30	411	40	160	66	328	40	315	103	309	115
291	77	6	268	34	279	93	287	35	329	3	294	93	305	29	365	0	147	92
292	222	61	197	10	2	10	41	26	180	108	229	100	199	94	71	83	542	113
293	39	44	45	24	117	30	60	24	199	53	74	85	58	54	438	9	274	6
294			261	1	134	11	231	117	194	17	299	27	419	81	225	50	526	34
295			161	47	218	72	58	100	153	42	391	35	22	119	189	74	166	34
296			314	116	344	109	190	33	92	59	244	48	293	100	418	68	57	88
297			234	96	26	39	164	18	94	110	132	33	285	35	306	118	498	101
298			249	67	44	98	1	42	6	115	143	21	207	65	335	37	502	10
299			79	49	267	66	95	44	414	15	390	94	475	113	152	39	254	109
300			5	75	53	47	377	55	192	73	179	12	360	56	472	72	422	8
301			220	42	148	61	99	74	46	86	59	40	25	22	463	53	108	8
302			264	52	237	86	176	67	352	60	298	110	284	8	136	109	85	17
303			49	30	79	69	18	72	63	25	315	63	74	28	186	47	81	0
304			80	39	82	13	25	78	401	11	207	46	290	114	70	80	226	44
305			245	103	337	114	284	21	321	13	409	67	66	20	130	60	35	26
306			11	107	309	33	342	114	370	95	197	83	252	96	101	6	160	106
307			287	17	244	29	365	16	189	94	431	70	69	47	404	78	80	10
308			50	7	45	104	339	47	111	71	357	102	48	104	104	51	194	18
309			18	81	225	39	271	53	80	39	54	52	15	42	53	4	385	85
310			106	23	110	116	94	68	168	21	248	49	340	44	285	91	381	68
311			295	75	157	79	219	103	304	26	215	34	34	71	10	16	105	5
312			291	8	31	25	325	90	381	53	89	32	243	53	317	20	351	23
313			186	38	220	113	12	29	279	80	335	28	369	86	235	98	7	14
314			35	62	234	55	299	84	135	18	60	51	215	2	360	6	215	13
315			16	27	41	5	326	97	114	26	113	116	467	59	241	1	452	99
316			178	104	351	102	300	39	55	3	342	47	355	1	282	34	317	20
317			205	12	268	17	362	83	202	2	334	107	401	55	430	8	9	67
318			300	55	0	12	0	55	314	111	417	55	390	17	79	19	234	16
319			46	61	203	16	282	80	223	106	356	61	178	68	331	52	192	21
320			311	56	95	89	124	13	178	78	84	114	424	49	73	101	138	78
321			251	118	212	5	221	58	1	74	24	5	197	103	64	30	295	36
322			211	40	222	105	142	98	160	81	287	119	117	43	231	5	214	3

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323			164	97	199	35	132	4	289	59	252	44	347	34	288	73	395	60
324					137	9	140	59	182	104	418	30	376	15	138	117	43	82
325					194	41	166	2	404	6	260	4	306	116	131	68	83	93
326					336	82	291	49	54	97	33	81	39	24	491	40	515	93
327					276	107	66	15	90	95	112	41	186	5	307	95	157	65
328					147	91	197	79	386	48	159	80	373	33	77	71	350	118
329					266	68	144	106	403	27	205	7	60	118	502	101	380	76
330					58	4	375	0	113	62	226	68	239	75	128	106	481	42
331					57	75	301	100	132	85	393	64	65	6	22	107	51	29
332					310	115	62	39	71	30	341	6	171	108	146	108	321	91
333					200	38	17	76	126	119	128	113	385	87	9	31	235	73
334					63	81	55	37	384	12	247	59	316	62	271	33	342	98
335					228	110	246	26	342	109	19	9	323	58	439	46	249	80
336					261	37	385	61	278	4	317	20	44	117	122	99	311	62
337					353	117	309	101	96	24	231	101	77	36	240	36	345	7
338					100	50	107	0	256	31	243	40	118	77	310	59	348	55
339					103	93	38	33	3	117	157	112	26	65	118	81	256	45
340					287	59	323	44	398	0	286	54	257	119	395	115	328	52
341					59	42	178	64	83	68	339	62	21	94	509	35	96	113
342					166	57	170	93	24	46	348	18	98	92	216	53	267	97
343					159	40	308	17	390	83	78	93	213	11	444	75	18	45
344					60	21	146	110	229	75	239	82	359	99	19	45	91	74
345					20	58	145	30	339	43	400	1	231	0	13	77	111	99
346					42	36	11	58	157	70	127	77	309	85	452	84	499	19
347					114	66	233	9	144	45	382	115	99	46	424	74	463	59
348					256	42	222	96	259	20	254	22	47	69	250	89	451	30
349					32	15	153	77	101	29	55	111	242	101	272	65	140	108
350					102	56	237	22	167	92	293	106	299	36	446	114	313	116
351					243	13	248	66	350	17	152	105	161	84	325	85	536	106
352					153	32	182	71	228	118	11	28	212	16	290	72	523	91
353					5	102	100	94	337	87	0	61	12	112	111	64	59	37
354							92	91	346	74	189	29	37	33	42	3	525	22
355							354	87	140	47	419	71	286	47	5	27	429	56
356							165	114	244	10	253	118	358	14	332	15	73	62
357							364	105	35	52	445	82	407	24	427	1	308	54
358							243	14	380	68	96	60	461	8	47	87	79	98
359							320	102	38	90	329	111	75	93	259	2	477	65
360							175	82	340	104	141	19	468	97	476	99	532	79
361							356	59	305	36	164	47	70	79	421	109	82	69
362							44	9	230	40	2	26	283	22	114	44	40	70
363							261	85	331	47	372	42	214	18	3	15	360	74
364							321	91	165	41	25	12	100	97	431	17	69	28
365							355	6	263	35	171	67	3	30	208	70	466	18
366							196	27	40	103	340	16	262	78	495	47	132	77
367							31	118	128	60	235	43	219	95	494	77	409	24
368							241	31	77	43	432	91	87	8	246	86	287	117
369							305	13	84	15	91	91	301	105	67	40	243	100
370							188	40	51	116	217	98	41	3	369	82	412	33
371							314	76	216	48	380	9	1	41	211	42	487	95
372							113	83	328	25	92	45	363	76	109	14	113	42
373							104	112	166	61	58	35	445	63	403	50	290	11
374							9	89	251	77	138	94	462	4	457	7	493	45
375							262	115	17	66	22	73	126	81	59	66	37	117
376							33	52	75	14	81	18	201	59	125	90	286	87
377							192	72	15	111	188	103	319	115	153	67	109	59
378							360	68	32	79	258	76	416	9	367	117	388	49
379							13	87	149	84	285	90	238	74	113	62	115	33
380							138	42	56	32	174	65	411	65	382	100	134	104
381							52	111	271	115	98	16	375	117	347	24	340	78
382							125	107	369	94	345	58	418	108	168	55	354	114
383							194	90	409	80	31	108	55	44	57	7	49	85
384							162	63	139	38	370	21	268	98	54	92	495	46
385							155	101	195	58	6	34	430	100	141	26	459	83
386							109	43	181	8	249	72	321	110	233	43	231	10
387									117	50	190	96	313	107	135	49	127	103
388									225	22	199	109	477	57	498	22	199	53

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389									118	55	318	117	320	104	468	37	244	2
390									112	107	236	32	349	54	456	113	430	47
391									336	62	398	87	297	38	295	111	212	57
392									375	66	71	89	258	27	65	105	14	13
393									284	101	304	13	223	37	311	66	490	14
394									288	21	44	46	209	32	467	27	171	75
395									236	96	381	104	417	39	11	43	472	119
396									296	57	90	37	180	17	255	112	533	29
397									82	102	352	89	5	72	297	63	151	110
398									224	65	43	104	53	70	182	29	167	32
399									276	9	52	69	114	34	386	116	319	35
400									11	89	222	36	255	63	263	84	266	38
401									345	112	374	48	458	42	126	12	124	25
402									270	69	192	38	62	99	393	3	324	39
403									30	34	240	84	264	13	286	68	31	44
404									47	6	355	66	281	66	479	71	418	111
405									376	61	110	99	144	82	215	16	440	36
406									306	32	367	5	247	28	465	5	494	80
407									365	38	76	50	314	20	497	56	101	41
408									214	19	336	105	120	26	461	92	200	115
409									155	25	288	30	205	33	252	78	25	58
410									185	87	330	81	191	25	253	113	371	31
411									235	91	449	79	122	89	27	96	298	4
412									408	105	265	92	429	16	492	49	164	95
413									20	57	436	78	152	29	484	67	283	50
414									164	112	233	33	453	52	334	45	316	72
415									85	33	415	112	97	86	477	51	476	71
416									326	85	63	30	139	21	312	79	142	79
417											166	31	174	41	407	119	483	73
418											227	116	79	116	236	115	71	27
419											448	118	365	50	194	11	177	48
420											122	63	282	57	96	1	255	119
421											399	10	414	114	106	46	497	52
422											257	102	217	31	139	91	275	0
423											325	24	372	45	405	96	518	76
424											42	58	277	55	242	25	365	67
425											309	0	182	118	75	54	524	48
426											208	17	7	1	398	13	333	29
427											162	20	102	12	390	107	32	97
428											65	59	273	6	299	9	173	26
429											272	75	9	68	25	59	89	101
430											216	15	260	10	371	31	11	40
431											39	74	46	49	412	62	204	51
432											433	14	232	92	480	57	312	5
433											140	48	162	54	496	69	543	114
434											274	85	450	75	219	118	428	102
435											368	40	367	5	33	17	56	38
436											66	66	45	45	247	22	373	32
437											213	86	198	61	507	110	175	102
438											148	1	295	69	239	83	88	22
439											323	34	288	4	30	84	239	23
440											328	45	378	18	144	97	280	64
441											360	93	124	2	356	90	535	108
442											291	57	112	19	204	94	54	52
443											107	7	228	73	249	114	225	1
444											17	103	57	50	345	48	181	68
445											283	99	394	88	353	102	90	100
446											278	94	188	91	485	33	270	92
447											15	14	474	36	160	31	520	89
448											124	5	85	103	56	13	344	40
449											347	49	204	35	373	55	119	46
450													438	79	172	21	368	43
451													23	55	88	79	273	81
452													13	102	48	89	197	62
453													196	109	273	110	213	98
454													412	27	20	33	341	110

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455																476	60	183	119	293	17
456																59	90	414	11	257	58
457																341	94	448	38	441	6
458																366	82	433	28	189	9
459																322	70	378	104	401	27
460																312	97	401	76	488	38
461																362	25	155	98	188	105
462																454	51	228	3	61	66
463																146	20	419	73	6	71
464																86	95	171	69	356	45
465																96	16	15	99	47	114
466																165	88	199	4	288	101
467																172	62	499	111	379	28
468																479	26	116	95	10	5
469																115	65	55	53	36	111
470																251	87	117	9	506	60
471																459	10	276	87	534	54
472																473	31	132	56	264	86
473																128	101	133	102	352	12
474																226	1	441	20	128	82
475																131	0	451	34	16	51
476																156	30	466	70	416	77
477																352	119	333	75	425	104
478																210	8	293	23	468	58
479																317	86	177	41	433	112
480																326	7	336	53	129	15
481																52	31	173	6	529	55
482																145	53	510	117	217	85
483																		383	2	279	47
484																		298	7	58	34
485																		258	103	155	93
486																		99	18	297	69
487																		279	67	318	109
488																		123	107	501	87
489																		489	82	411	23
490																		316	46	470	66
491																		112	0	144	3
492																		0	101	13	91
493																		149	76	74	112
494																		281	60	474	118
495																		43	81	530	75
496																		260	57	545	101
497																		512	79	531	94
498																		432	36	107	31
499																		435	35	137	7
500																		321	39	165	70
501																		244	62	60	43
502																		76	42	208	76
503																		305	19	359	79
504																		167	74	98	107
505																		179	105	78	2
506																		185	27	121	45
507																		506	18	187	116
508																		91	100	482	86
509																		420	15	358	3
510																		470	91	426	85
511																		328	12	72	81
512																		511	50	210	54
513																				46	19
514																			250	118	
515																				507	92
516																				454	117
517																				438	98
518																				540	11
519																				376	26
520																				325	115

EXPERIMENTAL SPECIFICATION FOR FLEXIBLE SCC TURBO CODE FOR TELEMETRY APPLICATIONS

521																	183	73
522																	158	96
523																	294	83
524																	402	107
525																	377	42
526																	383	22
527																	460	44
528																	449	13
529																	176	21
530																	334	59
531																	63	83
532																	241	56
533																	170	87
534																	327	111
535																	276	63
536																	149	72
537																	163	71
538																	413	8
539																	260	67
540																	419	72
541																	420	14
542																	114	8
543																	434	39
544																	306	103
545																	432	21

ANNEX D

PARALLELISED INTERLEAVER DESIGN

The methodology used to design good interleavers is based on the classical approach of random incremental generation of the permutation law and application of spread constraints, as described below:

The spread constraint was the following (linear):

$$\frac{|i-j|}{S_1} + \frac{|\pi(i) - \pi(j)|}{S_2} \geq 1 \quad (\text{B-1})$$

The general permutation law π satisfying the parallel access constraints is described through the two vectors β and α describing the shift and macro address permutation as follows:

$$\pi(i) = W \left\{ \left\lfloor i/W \right\rfloor + \beta(i_w) \right\}_{120} + \alpha(i_w) \quad i = 0, \dots, N-1 \quad (\text{B-2})$$

where the notation $\{\}_{N}$ denotes the remainder or modulus operation, and $W=N/120$ is the size of each memory bank.

The algorithm works as follows:

- 1) Fix the value of the parameters $S_1=S_2=\sqrt{N}$, which is the theoretical maximum one. Fix the maximum number of local trials T_l and the maximum number of global trials T_g .
- 2) Sequentially, for each possible i ranging from 0 to $W-1$.
- 3) Generate a new random value of $\beta(i)$ in the range $[0,119]$ and a new random value for the permutation $\alpha(i)$ on W elements.
- 4) The specification of these new values allows to be determined a subset of generally non-contiguous values of the permutation $\tilde{\pi}$.
- 5) Check the condition (B-1) on all positions j for which the permutation is already defined.
- 6) If the condition is verified continue to next i from point 2; otherwise increase the number of local trials.
- 7) If the number of local trials is smaller than T_l go back to point 3; otherwise increase the number of global trials.
- 8) If the number of global trials is smaller than T_g restart from point 2, setting $i=0$.
- 9) Alternatively reduce S_1 and S_2 by one and restart.

This procedure allows good interleavers to be obtained in an amount of time that can be controlled through the choice of T_l and T_g . The upper row of the tables describing the interleavers report the obtained values of S_1 and S_2 .

ANNEX E

PHYSICAL LAYER PSEUDO-RANDOMIZATION

The use of Physical layer pseudo-randomization is envisaged as an option for energy dispersal. The PL pseudo-randomization is obtained by multiplying the In-phase and Quadrature samples by a complex randomization sequence ($C_I + jC_Q$):

$$I_{\text{randomized}} = \{ I \cdot C_I - Q \cdot C_Q \} \qquad Q_{\text{randomized}} = \{ Q \cdot C_I + I \cdot C_Q \}$$

The randomization sequence is not applied to the PL header. The randomization sequence is reinitialized at the end of each Physical layer header (FM+FD).

The scrambling code sequences are constructed by combining two real m -sequences (generated by means of two generator polynomials of degree 18) into a complex sequence. The resulting sequences thus constitute segments of a set of Gold sequences. Let x and y be the two sequences respectively.

The x sequence is constructed using the primitive polynomial $h(x) = 1 + x^7 + x^{18}$.

The y sequence is constructed using the polynomial: $g(y) = 1 + y^5 + y^7 + y^{10} + y^{18}$.

The sequence depending on the chosen scrambling code number n is denoted z_n in the sequel. Furthermore, let $x(i)$, $y(i)$ and $z_n(i)$ denote the i^{th} symbol of the sequence x , y , and z_n respectively.

The m -sequences x and y are constructed as:

- Initial conditions:
 - x is constructed with $x(0) = 1, x(1) = x(2) = \dots = x(16) = x(17) = 0$.
 - $y(0) = y(1) = \dots = y(16) = y(17) = 1$.
- Recursive definition of subsequent symbols:
 - $x(i+18) = x(i+7) + x(i) \text{ modulo } 2, i = 0, \dots, 2^{18} - 20$.
 - $y(i+18) = y(i+10) + y(i+7) + y(i+5) + y(i) \text{ modulo } 2, i = 0, \dots, 2^{18} - 20$.
- The n^{th} Gold code sequence $z_n, n = 0, 1, 2, \dots, 2^{18} - 2$, is then defined as:
 - $z_n(i) = [x((i+n) \text{ modulo } (2^{18}-1)) + y(i)] \text{ modulo } 2, i = 0, \dots, 2^{18} - 2$.

These binary sequences are converted to integer valued sequences R_n (R_n assuming values 0, 1, 2, 3) by the following transformation:

$$R_n(i) = 2 z_n((i + 131\,072) \text{ modulo } (2^{18}-1)) + z_n(i), i = 0, 1, \dots, 133440.$$

Finally, the n^{th} complex scrambling code sequence $C_I(i) + jC_Q(i)$ is defined as:

$$C_I(i) + jC_Q(i) = \exp^{j R_n(i) \pi/2}$$

R_n	$\exp^{j R_n \pi/2}$	$I_{\text{randomized}}$	$Q_{\text{randomized}}$
0	1	I	Q
1	j	$-Q$	I
2	-1	$-I$	$-Q$
3	$-j$	Q	$-I$

The figure below gives a possible block diagram for pseudo-randomization sequence generation for $n = 0$.

